



LEGALM: Efficient Legalization for Mixed-Cell-Height Circuits with Linearized Augmented Lagrangian Method

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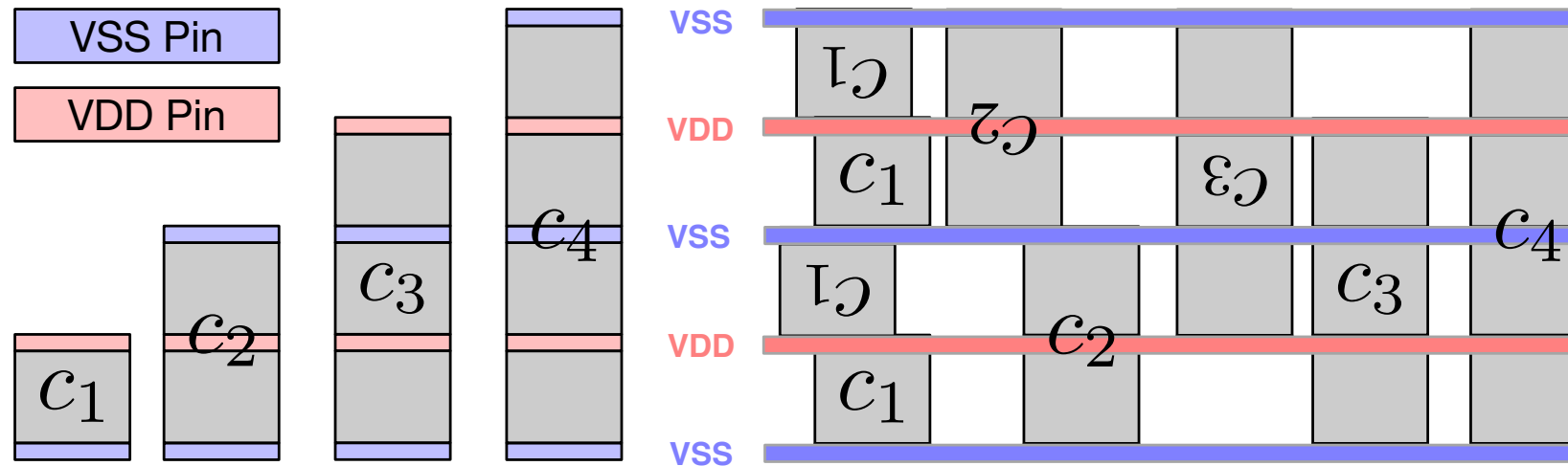
1. Introduction
2. The LEGALM Algorithm
3. Experimental Results
4. Conclusion

Introduction

Mixed-Cell-Height Circuits



- Recent mixed-cell-height designs combine higher and smaller cells to optimize PPA in modern ASICs.
 - Higher cells boost performance and routability for critical paths.
 - Smaller cells improve area efficiency and reduce power for non-critical logic.

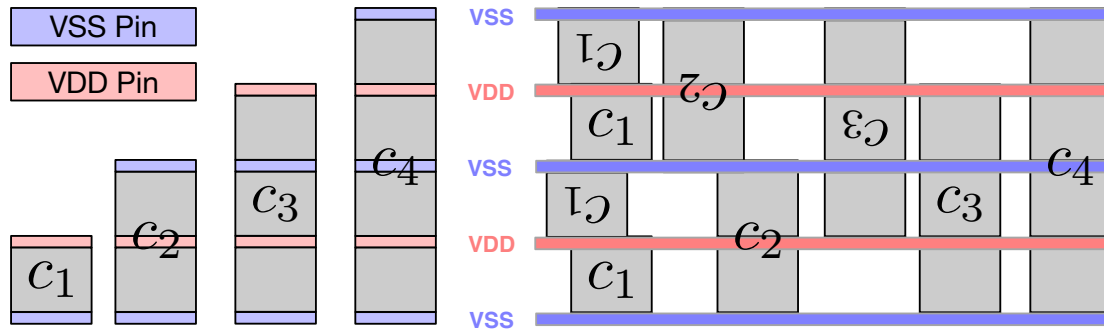


Challenges & Motivation



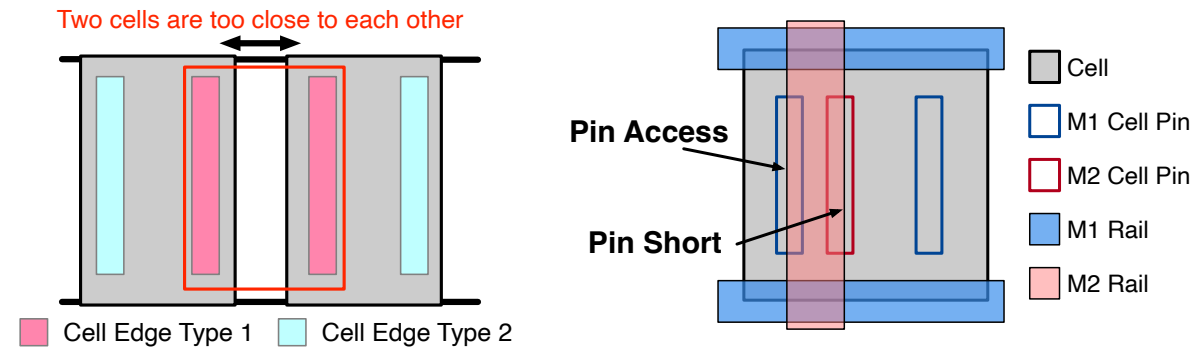
Why Legalization matters:

- ▶ Eliminates design rule violations post-global placement.
- ▶ Impacts downstream routing and performance.



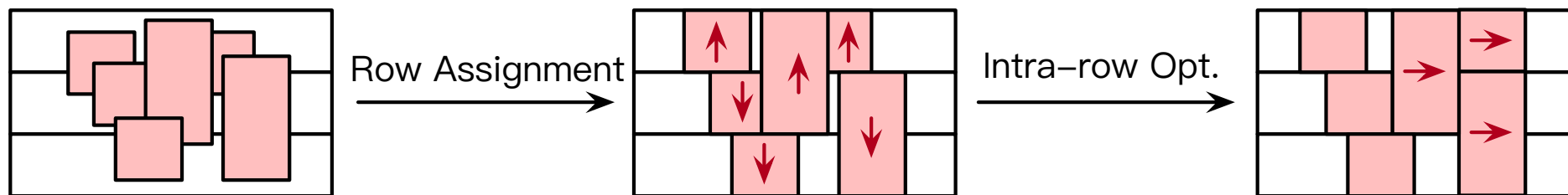
Challenges in Mixed-Cell-Height:

- ▶ Site alignment, overlap-free.
- ▶ Cross-row shapes, P/G alignment, fence regions, edge spacing, pin access.



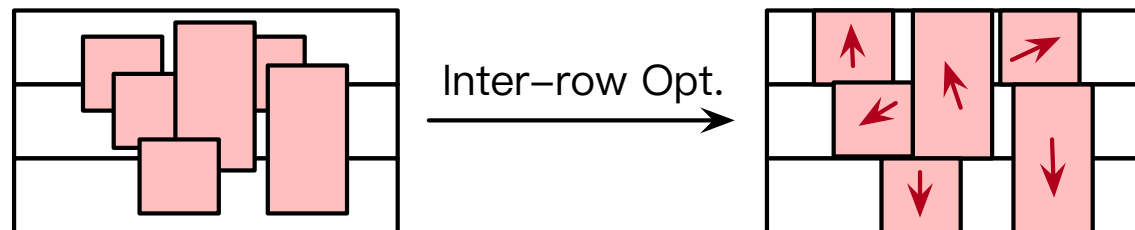
□ Intra-row Methods:

- ▶ Row assignment + intra-row optimization algorithm (e.g., ILP [Li+,DAC'18], LCP [Chen+, DAC'17]).
- ▶ Limited vertical movement space during optimization.



□ Inter-row Methods:

- ▶ Network flow / ILP-based (e.g., [Darav+, ISPD'17])
- ▶ High flexibility but computationally expensive.



Gap:

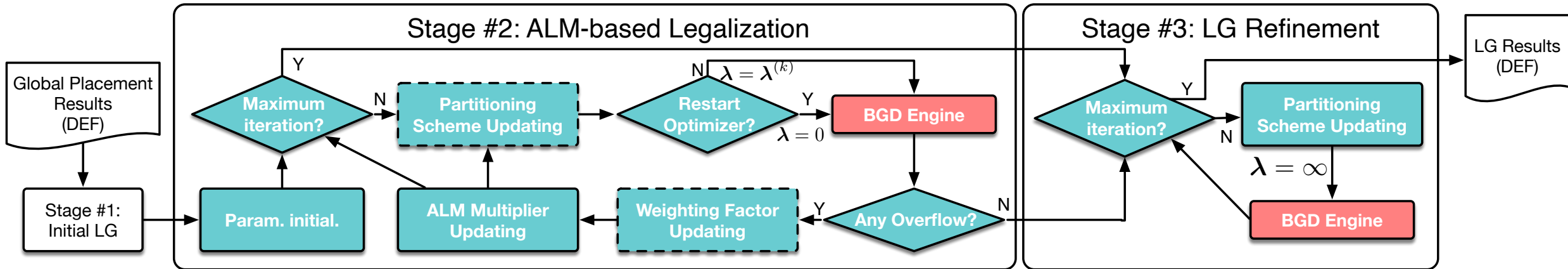
Existing methods struggle with complex constraints and scalability.

LEGALM, an efficient legalization method for mixed-cell-height circuits using the linearized augmented Lagrangian method.

1. **Augmented Lagrangian Method (ALM)** for efficient vertical and horizontal cell movements.
2. **Block Gradient Descent (BGD)** for parallel cell updates.
3. **Triple-fold partitioning** for GPU efficiency. $94.2 \times$ speedup with less than 0.5% quality degradation.

Results: 6-36% better quality, 2.25 - $5.99 \times$ speedup on million-cell designs.

The LEGALM Algorithm



❑ 3-Stage Workflow:

1. **Initial Legalization:** Minimal displacement, ignore overlaps.
2. **ALM-based Legalization:** Optimize displacement + eliminate overlaps.
3. **Refinement:** Strict no-overlap optimization.

❑ Key Components:

- Linearized ALM formulation for constraint relaxation.
- BGD for parallel cell updates.
- GPU-friendly partitioning.

Augmented Lagrangian Formulation



- Objective: Minimize displacement while satisfying constraints.

$$\min \sum_{j \in S} \sum_{i \in N} \sum_{t \in \mathcal{T}_i} w_{i,t,j} x_{i,t,j},$$

$$\text{s.t. } x_{i,t,j} \in \{0, 1\},$$

$$g_j(\mathbf{x}) = \sum_{i \in N} \sum_{t \in \mathcal{T}_i} x_{i,t,j} - 1 \leq 0, \quad \forall j \in S,$$

connected sub-cell constraints,

fence region constraints,

- \mathcal{T}_i Partition cell i into sub-cells of width 1 and height H (row height).
- $x_{i,t,j}$ Binary indicator if sub-cell t ($t \in \mathcal{T}_i$) of cell i is placed at site j ($j \in S$).
- $w_{i,t,j}$ Displacement cost for placing sub-cell t ($t \in \mathcal{T}_i$) at site j ($j \in S$).
- $g_j(\mathbf{x})$ Overflow at site j ($j \in S$).

Augmented Lagrangian Formulation (cont'd)



- **Lagrangian relaxation** to handle overlap-free constraints.
- **Slack variables** r_j for site overflow.

$$\begin{aligned}\mathcal{L}(\mathbf{x}, \mathbf{r}, \boldsymbol{\lambda}) &= \sum_{j \in \mathcal{S}} \sum_{i \in \mathcal{N}} \sum_{t \in \mathcal{T}_i} w_{i,t,j} x_{i,t,j} \\ &\quad + \sum_{j \in \mathcal{S}} \lambda_j \left[(g_j(\mathbf{x}) + r_j) + \frac{\sigma}{2} (g_j(\mathbf{x}) + r_j)^2 \right] \\ &\quad + I_{\mathcal{X}}(\mathbf{x}),\end{aligned}$$

- Solve for optimal r_j using **elimination method**: $r_j = \max(0, -\frac{1}{\sigma} - g_j(\mathbf{x}))$.
- Iterative multiplier updates based on **KKT conditions**.

$$\mathbf{x}^{(k+1)} = \underset{\mathbf{x}}{\operatorname{argmin}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) = \underset{\mathbf{x}}{\operatorname{argmin}} \psi(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) + I_{\mathcal{X}}(\mathbf{x}).$$

$$\lambda_j^{(k+1)} = \max \left(\lambda_j^{(k)} + h_f \cdot \left[(g_j(\mathbf{x}) + r_j) + \frac{\sigma}{2} (g_j(\mathbf{x}) + r_j)^2 \right], 0 \right),$$

Linearized Proximal Gradient Method



$$\mathbf{x}^{(k+1)} = \underset{\mathbf{x}}{\operatorname{argmin}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) = \underset{\mathbf{x}}{\operatorname{argmin}} \psi(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) + I_{\mathcal{X}}(\mathbf{x}).$$

Iterative Optimization:

- In each iteration, given fixed $\boldsymbol{\lambda}$, let $f(\mathbf{x}) = \psi(\mathbf{x}, \boldsymbol{\lambda})$, solve the following function

$$\min_{\mathbf{x}} f(\mathbf{x}) + I_{\mathcal{X}}(\mathbf{x}).$$

- Proximal Mapping Operator (can be proved a linear operator).

$$P_{\mathcal{X}}(\mathbf{v}) = \arg \min_{\mathbf{u} \in \mathcal{X}} \|\mathbf{u} - \mathbf{v}\|_2^2,$$

- Proximal Gradient Method.

$$\mathbf{x}^{(k+1)} \in P_{\mathcal{X}}\left(\mathbf{x}^{(k)} - \tau \nabla f(\mathbf{x}^{(k)})\right),$$

Final Cost Function:

$$\text{cost}_{i,t,j} = w_{i,t,j} + \lambda_j \max\{1 + \sigma g_j(\mathbf{x}), 0\} + p \cdot H \cdot R_{i,t,j}$$

- $w_{i,t,j}$: displacement cost for placing sub-cell t ($t \in \mathcal{T}_i$) at site j ($j \in S$).
- $g_j(\mathbf{x})$: Overflow at site j ($j \in S$).
- $R_{i,t,j}$: total number of DRVs for placing sub-cell t ($t \in \mathcal{T}_i$) at site j ($j \in S$).

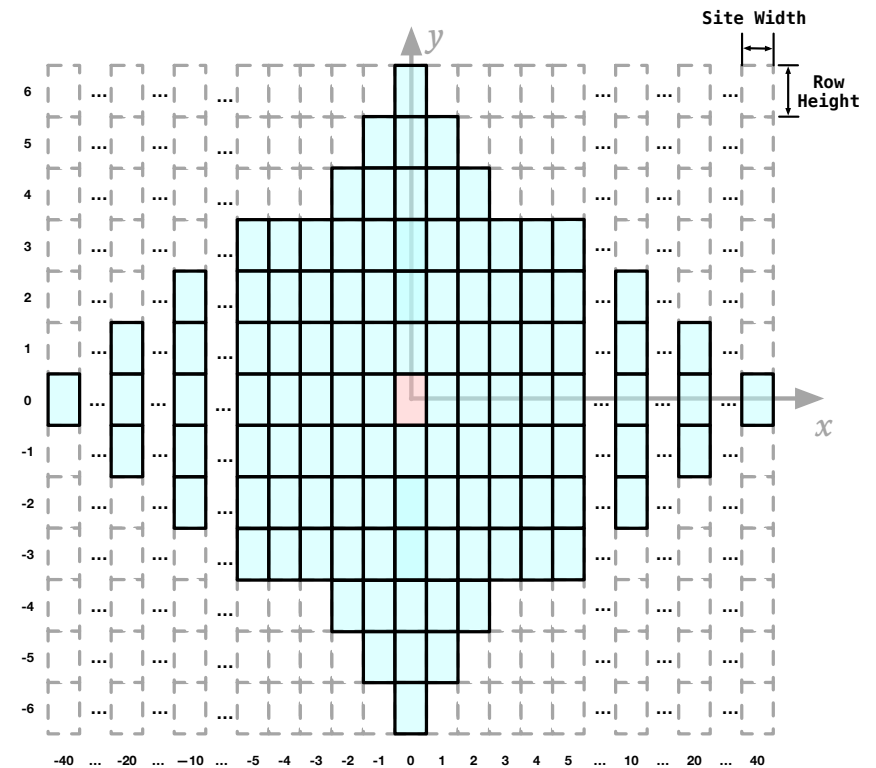


Figure: Diamond search range for the cell located at the red site.

Linearized Proximal Gradient Method (cont'd)



$$\mathbf{x}^{(k+1)} = \underset{\mathbf{x}}{\operatorname{argmin}} \mathcal{L}(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) = \underset{\mathbf{x}}{\operatorname{argmin}} \psi(\mathbf{x}, \boldsymbol{\lambda}^{(k)}) + I_{\mathcal{X}}(\mathbf{x}).$$

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Steps:

- Enumerate candidate positions (diamond search).
- Compute costs for displacement, overflow and technology.
- Select minimum-cost positions.

Advantages:

- Balance displacement and technology constraints.

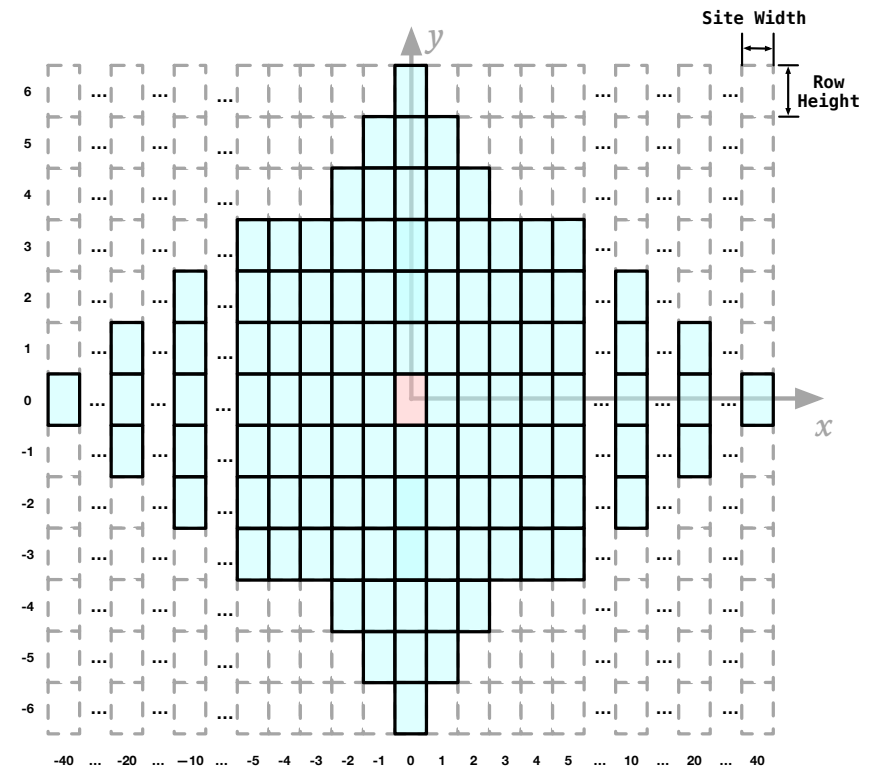
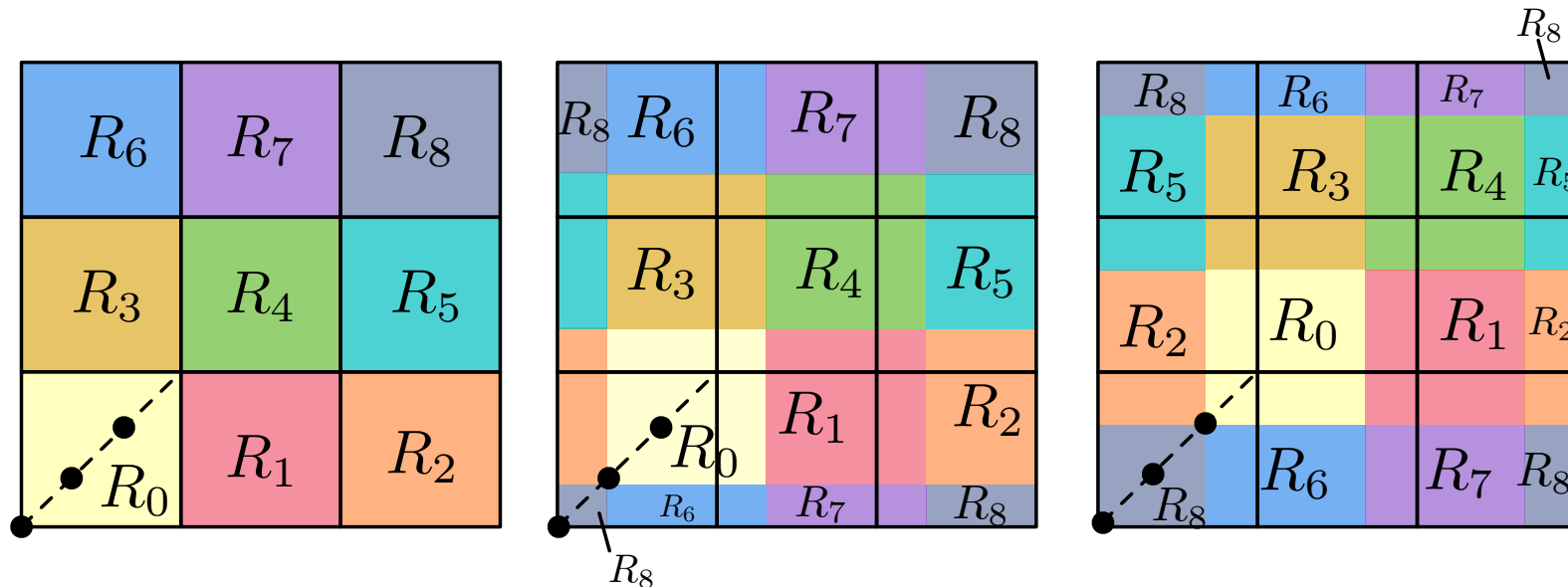
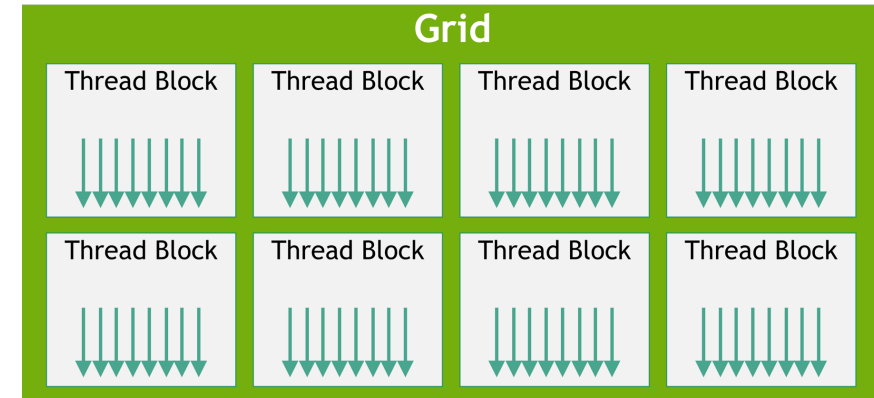


Figure: Diamond search range for the cell located at the red site.

Block Gradient Descent



- **Goal:** Parallelize cell updates without convergence loss.
- **Triple-fold Partitioning:**
 - Divide layout into non-overlapping sub-regions (totally 3 partition schemes).
 - Rotate partitions schemes to avoid update conflicts.
- **GPU Acceleration:**
 - Kernel design: 1 grid per sub region, threads for candidate selections.
 - Shared memory for cost reduction.



Experimental Results

Experiment Setup



□ Benchmarks:

- ICCAD-2017 (routability / fence regions). [Darav+, ICCAD'17]
- Modified ISPD-2015 (million-cell designs). [Chow+, DAC'16]

□ Metrics:

- Quality score \mathcal{S} : Combine HPWL variation \mathcal{S}_{hpwl} , weighted averaged displacement \mathcal{S}_{am} , maximum displacement \mathcal{M}_{max} , #pin short/access DRVs N_p , and #edge spacing DRVs N_e .
- Runtime and speedup.

□ Platform:

- One NVIDIA A800 GPU.
- Two Intel Xeon Platinum 8358 CPUs (2.60GHz, 32 cores) with 1024GB RAM.

Table: Statistics of ICCAD-2017 Benchmarks.

| Case | #Cells of Different Heights (H) | | | | Den. (%) | #Regions |
|--------------------|-------------------------------------|------|------|------|----------|----------|
| | 1 | 2 | 3 | 4 | | |
| des_perf_1 | 112644 | 0 | 0 | 0 | 90.6 | 0 |
| des_perf_a_md1 | 103589 | 4699 | 0 | 0 | 55.1 | 4 |
| des_perf_a_md2 | 105030 | 1086 | 1086 | 1086 | 55.9 | 4 |
| des_perf_b_md1 | 106782 | 5862 | 0 | 0 | 55.0 | 12 |
| des_perf_b_md2 | 101908 | 6781 | 2260 | 1695 | 64.7 | 12 |
| edit_dist_1_md1 | 118005 | 7994 | 2664 | 1998 | 67.4 | 0 |
| edit_dist_a_md2 | 115066 | 7799 | 2599 | 1949 | 59.4 | 1 |
| edit_dist_a_md3 | 119616 | 2599 | 2599 | 2599 | 57.2 | 1 |
| fft_2_md2 | 28930 | 2117 | 705 | 529 | 82.7 | 0 |
| fft_a_md2 | 27431 | 2018 | 672 | 504 | 32.3 | 0 |
| fft_a_md3 | 28609 | 672 | 672 | 672 | 31.2 | 0 |
| pci_bridge32_a_md1 | 26680 | 1792 | 597 | 448 | 49.5 | 4 |
| pci_bridge32_a_md2 | 25239 | 2090 | 1194 | 994 | 57.7 | 4 |
| pci_bridge32_b_md1 | 26134 | 1756 | 585 | 439 | 26.6 | 3 |
| pci_bridge32_b_md2 | 28038 | 292 | 292 | 292 | 18.3 | 3 |
| pci_bridge32_b_md3 | 27452 | 292 | 585 | 585 | 22.2 | 3 |

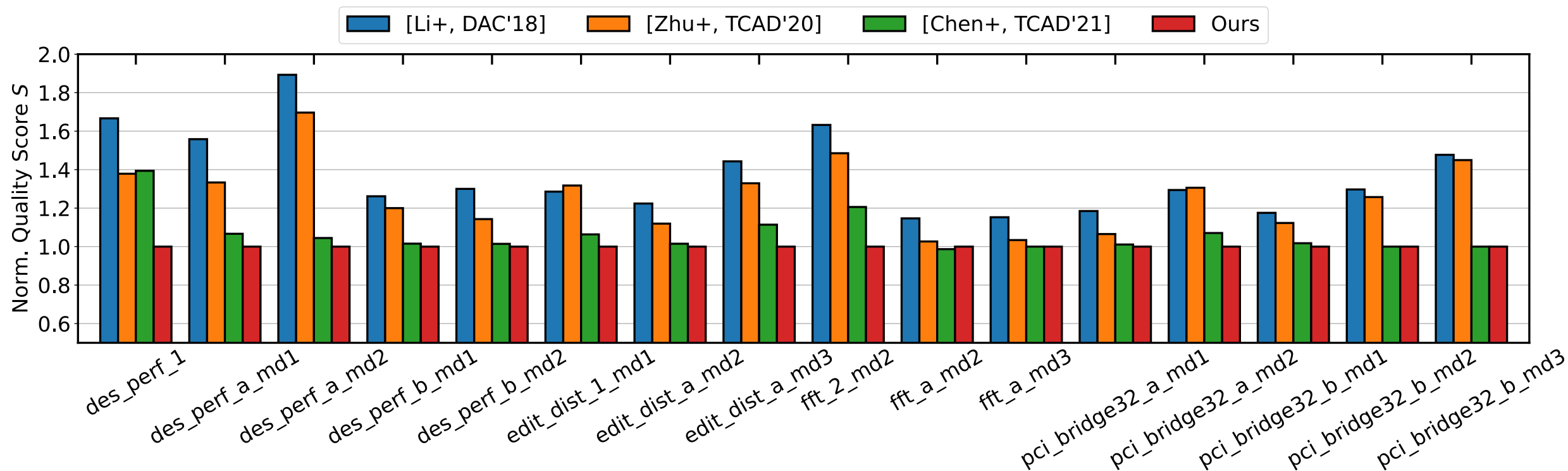
Table: Statistics of ISPD-2015 Benchmarks.

| Case | #Cells of Different Heights (H) | | Den. (%) |
|-------------------|-------------------------------------|--------|----------|
| | 1 | 2 | |
| mgc_superblue11_a | 861314 | 64302 | 43 |
| mgc_superblue12 | 1172586 | 114362 | 45 |
| mgc_superblue14 | 564769 | 47474 | 56 |
| mgc_superblue16_a | 625419 | 47474 | 48 |
| mgc_superblue19 | 478109 | 27988 | 52 |

Result on ICCAD-2017



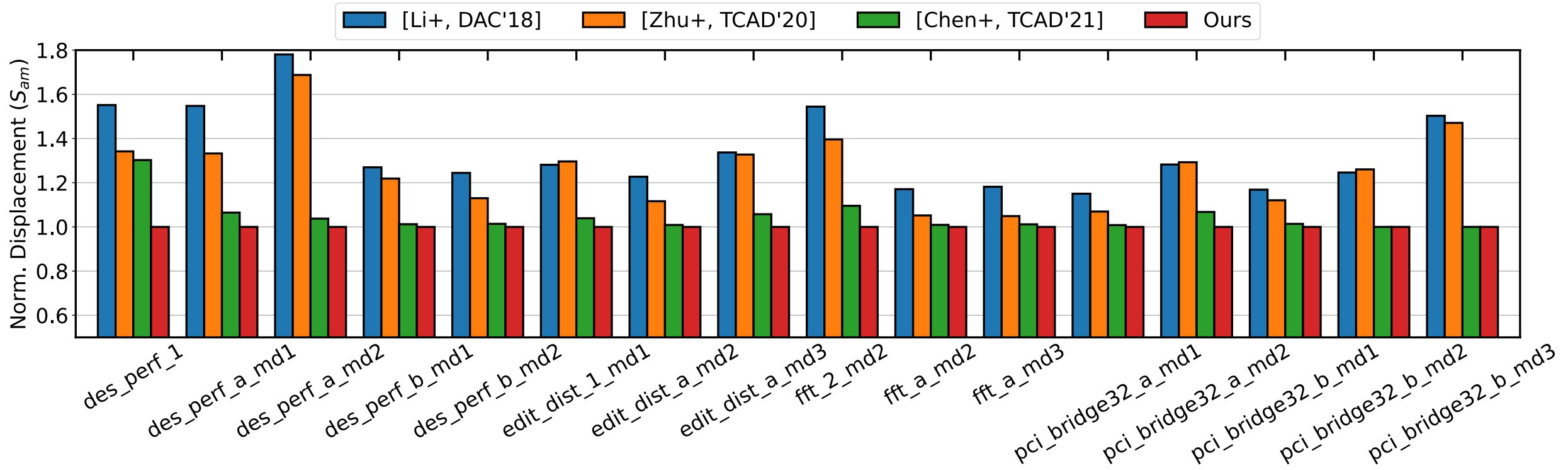
- Quality: LEGALM improve overall score S by 6-36% vs. SOTA with 1.03-3.83 \times speedup.



Result on ICCAD-2017 (cont'd)



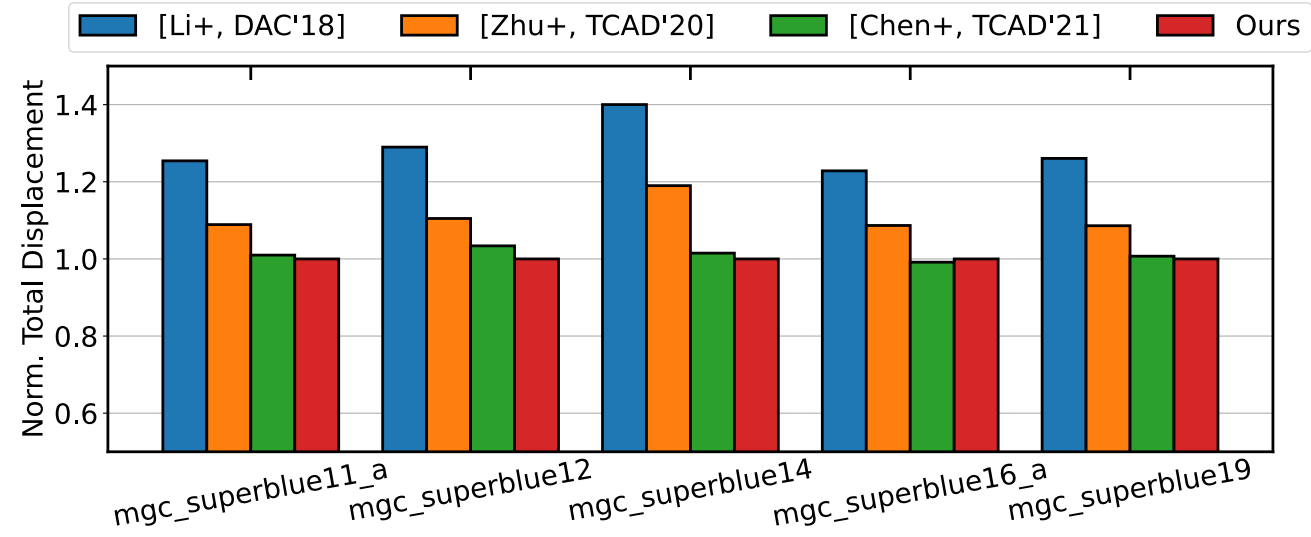
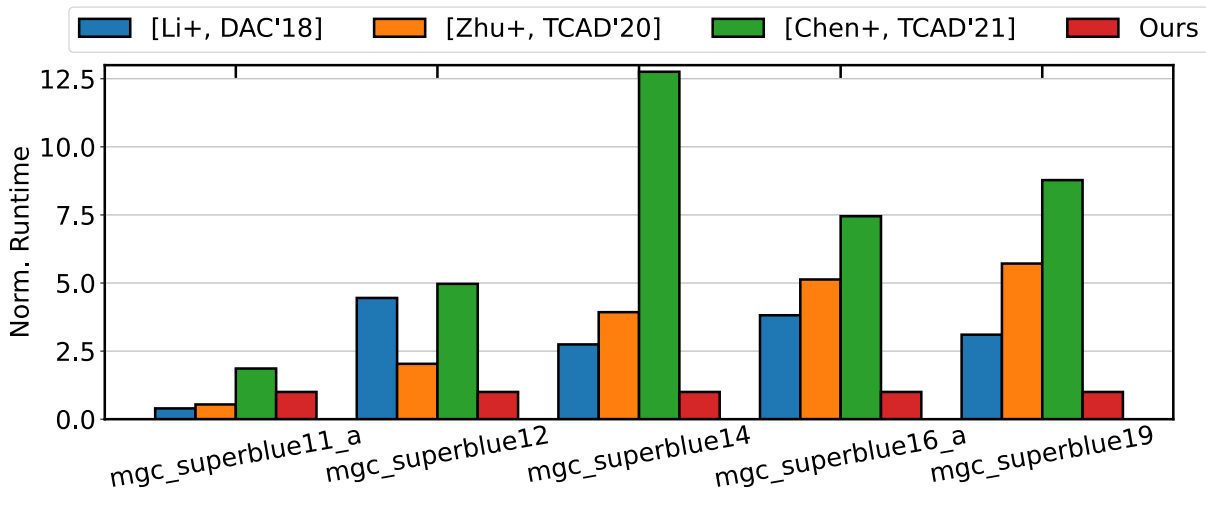
- Quality: LEGALM improve overall score S by 6-36% vs. SOTA with $1.03-3.83\times$ speedup.
- Breakdown:
 - 13-34% better HPWL variation S_{hpwl} .
 - 10% better maximum displacement M_{max} .
 - 4.4-33.2% better weighted average displacement S_{am} .



Scalability on Large Designs



- **Cases:** mgc_superblue (1M+ cells)
- **Results:**
 - 2.25–5.99× faster than SOTA.
 - 1.1–28.5% better displacement than SOTA.
 - Legalization in < 10 seconds for 3/5 cases.
- **Why?:** GPU parallelism + efficient partitioning.



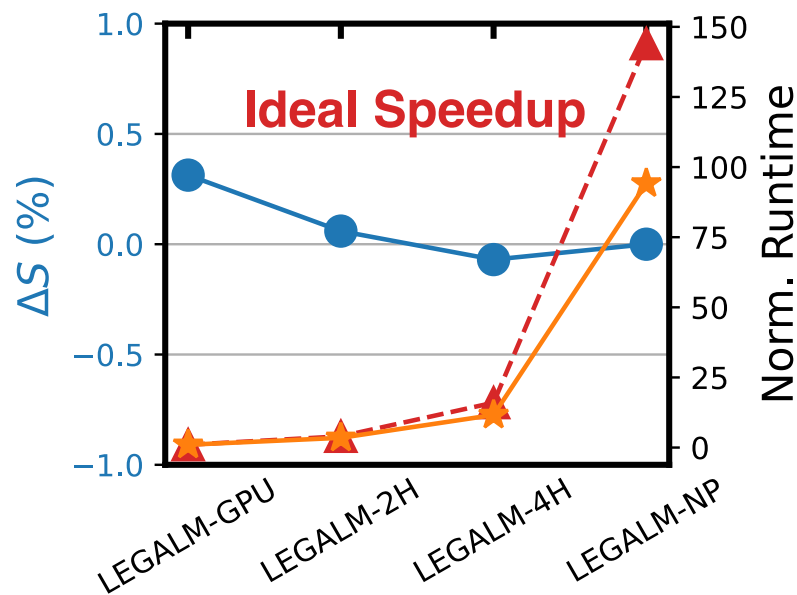
Ablation Study & Runtime Breakdown



Triple-fold Partitioning (TP) Impact:

- 94.2 \times speedup vs. no partitioning (LEGALM-NP).
- <0.5% quality loss.

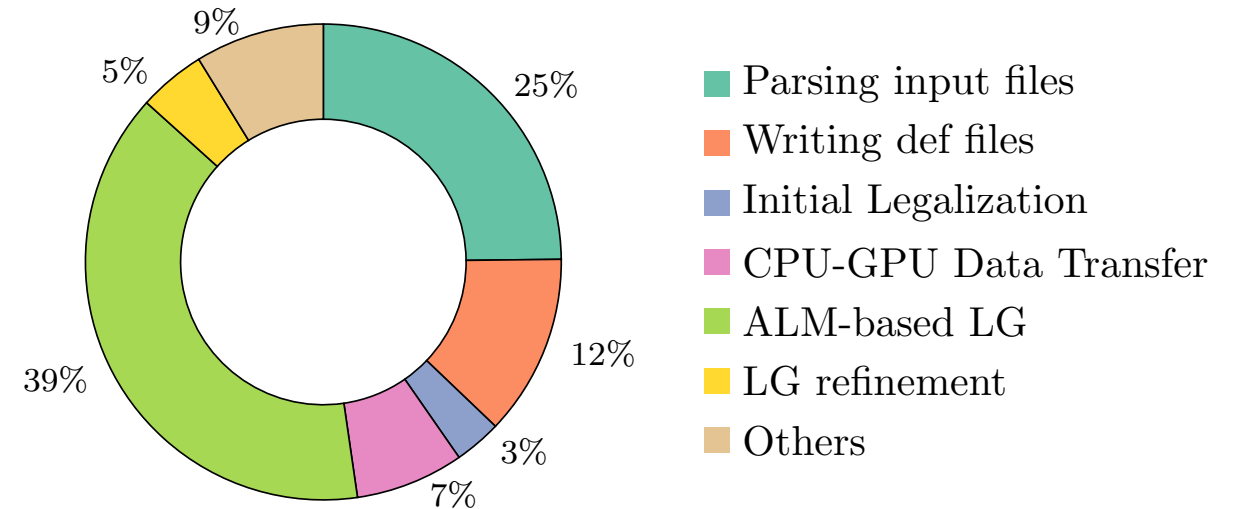
Grid Size: Larger sub-regions reduce parallelism but improve movement.



LEGALM-2H: double sub-region size
LEGALM-4H: quadruple sub-region size

Runtime Breakdown:

- ALM-based Legalization: 39%.
- Initial Legalization: 3%
- Legalization Refinement: 5%.



Conclusion & Future Work

LEGALM, an efficient legalization method for mixed-cell-height circuits using the linearized augmented Lagrangian method.

□ Conclusion:

- Linearized ALM formulation for mixed-cell-height legalization.
- Block Gradient Descent (BGD) + Triple-fold partitioning for GPU acceleration.
- 6–36% better quality, $94.2\times$ speedup.

□ Future Work:

- Advanced parallelization strategies.

THANK YOU!

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