

Robust Layer Assignment for Via Optimization in Multi-layer Global Routing

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Outline

Introduction



Motivation



Problem Formulation



Algorithm



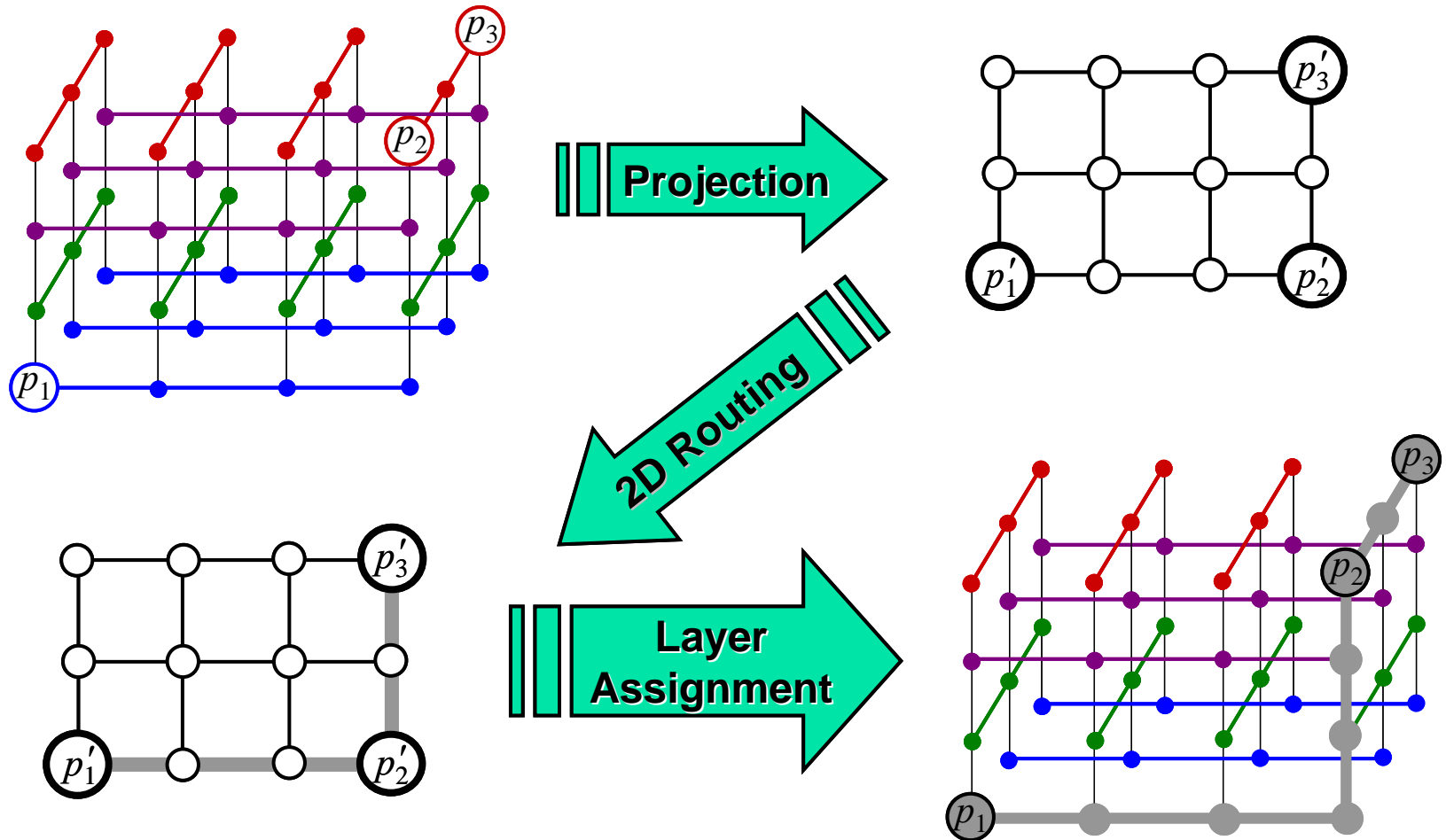
Experimental results



Conclusions

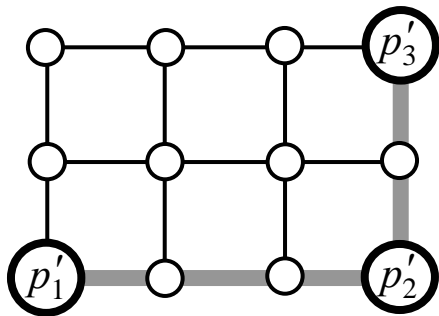
Introduction to Layer Assignment

- **Layer assignment** is a major step in multi-layer global routing

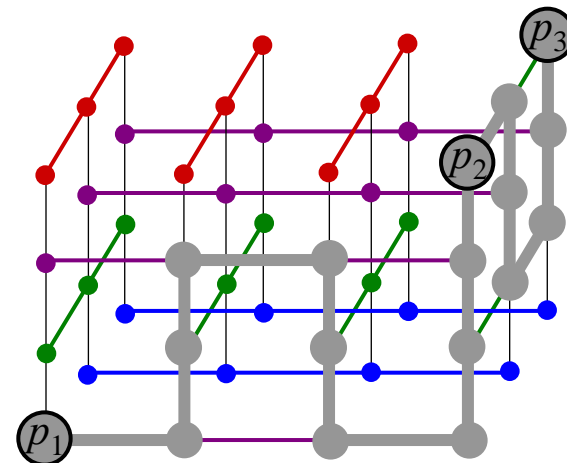


Introduction to Layer Assignment

- **Layer assignment** determines the final routing result
 - **A bad layer** assignment devastates all the previous efforts

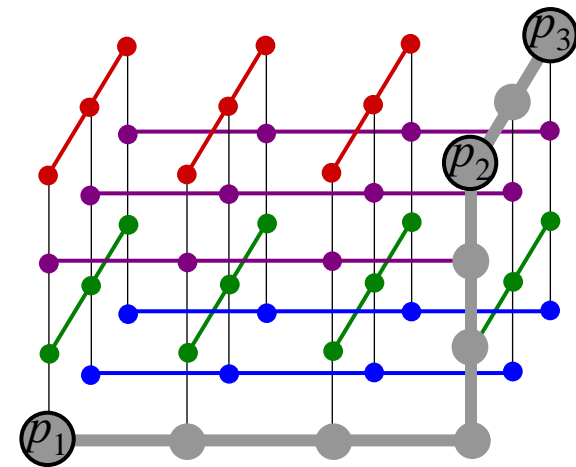
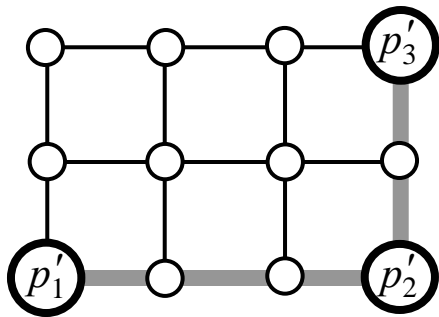


**Additional wire overflow
longer wirelength**



Introduction to Layer Assignment

- **Layer assignment** determines the final routing result
 - **A good layer** assignment keeps all the previous efforts

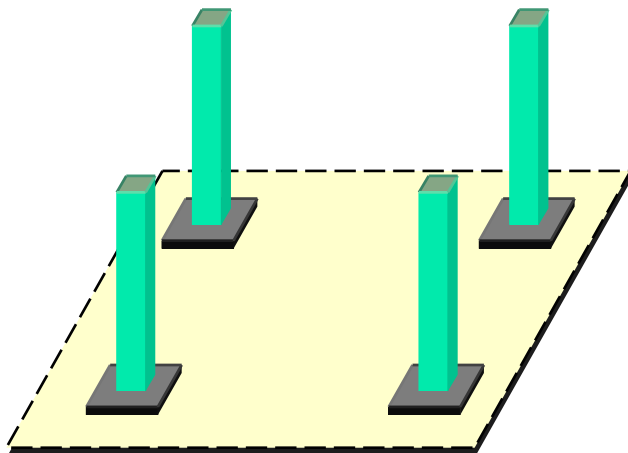


No additional wire overflow
Minimal wirelength

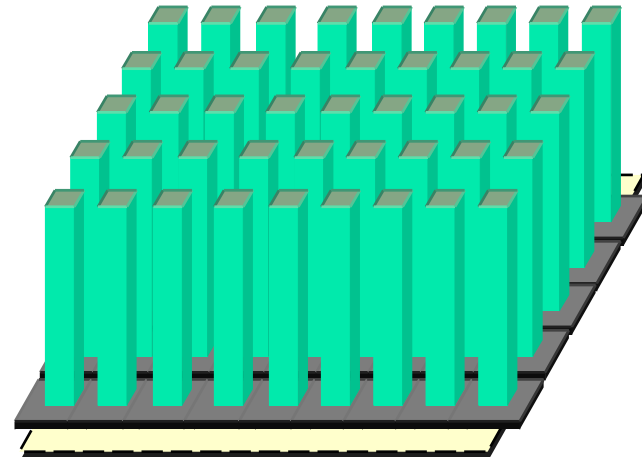


Motivation

- ❑ ISPD'07 and ISPD'08 Global Routing Contest did not limit # of **vias** placed in a tile



Allowable
in the contest

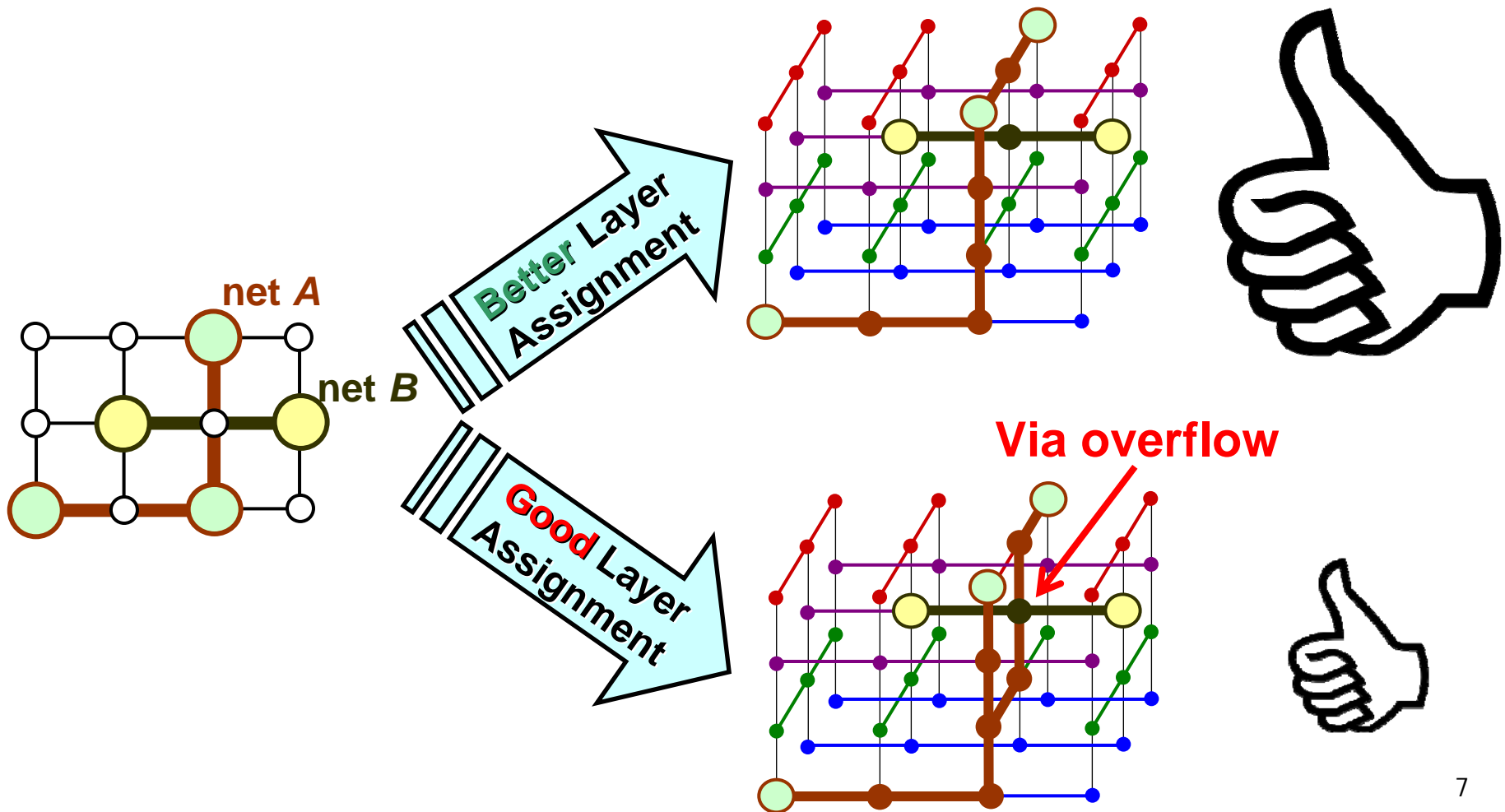


Still allowable
in the contest

- ❑ Routing result without considering **via capacity** is not practical !

Motivation

- A better layer assignment should take the **via capacity** into account



Previous Work for Via Capacity

□ [Hsu et al., ICCAD'08]

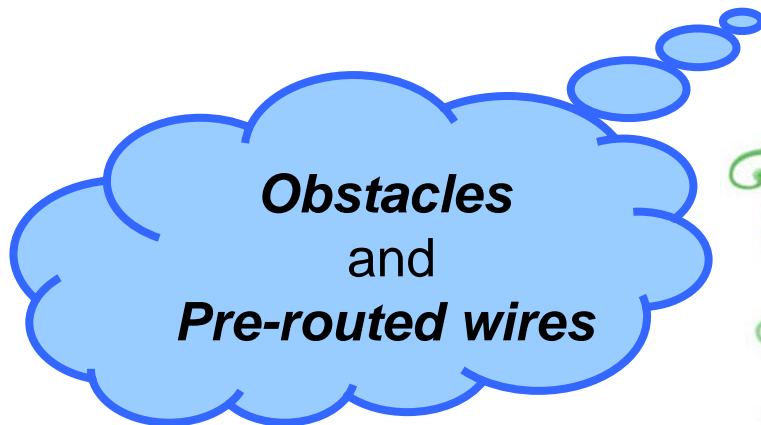
“Multi-layer Global Routing Considering Via and Wire Capacities”

- Considering via capacity for each tile
- No detailed information of its layer assignment step

□ Via capacity of a tile

= $remaining_area / via_area$

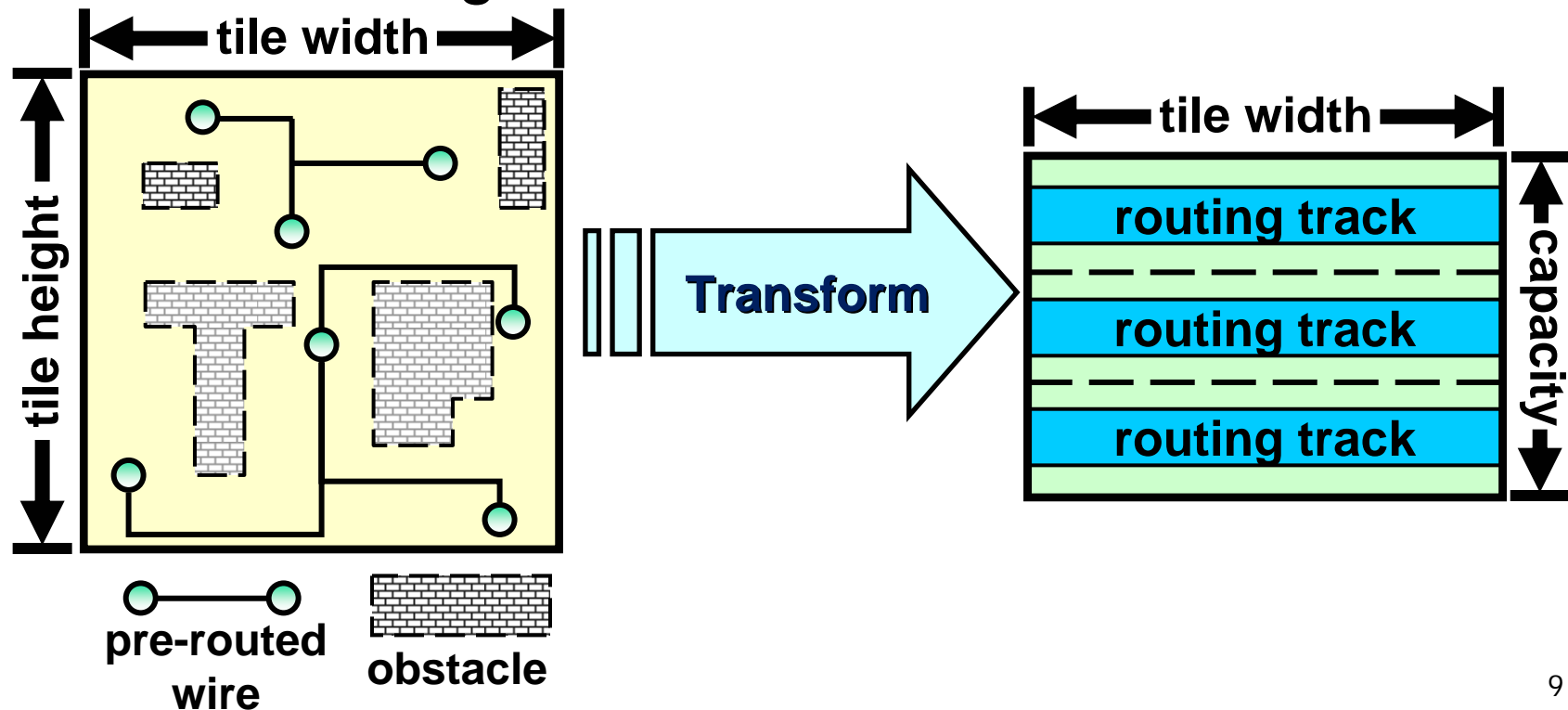
= $(tile_area - \underline{preoccupied_area}) / via_area$



**Didn't specify
in the benchmarks**

Problem Formulation

- ***“Congestion is modeled by including capacity adjustments. In the global routing benchmarks, there may be obstacles, or pre-routed wires.”*** — quoted from “details of file formats” of ISPD’08 Global Routing Contest rules

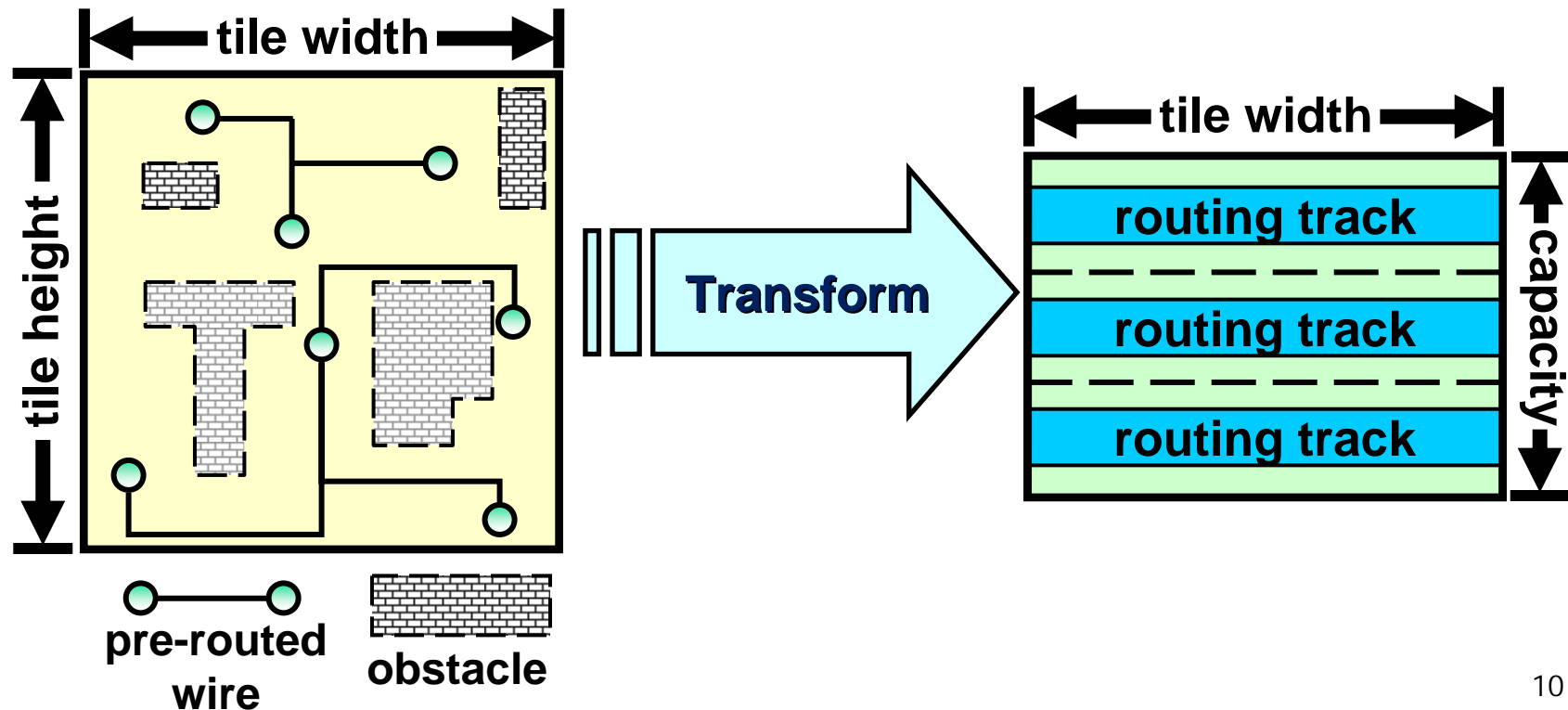


Problem Formulation

□ Via capacity of a tile

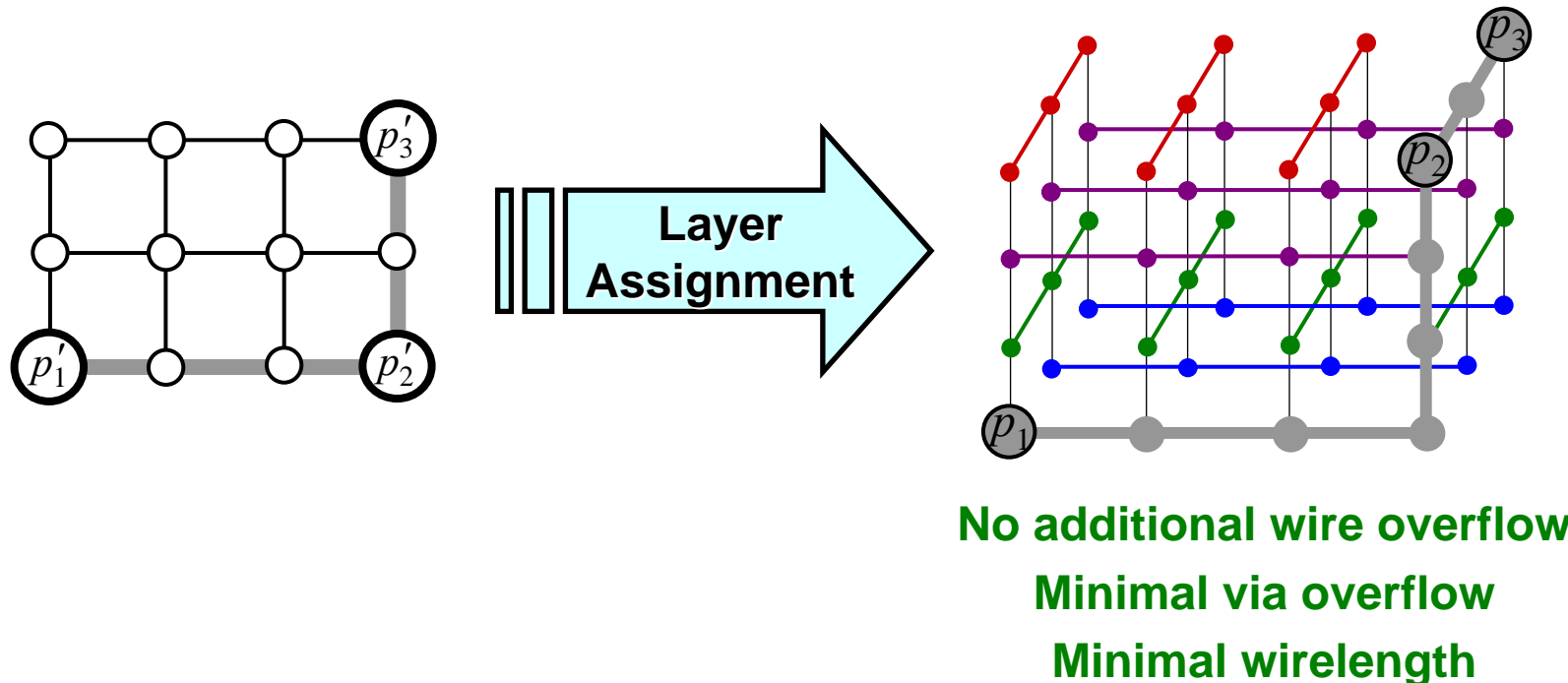
$$= (\text{tile_area} - \text{preoccupied_area}) / \text{via_area}$$

$$= (\text{capacity} \times \text{tile_width}) / \text{via_area}$$



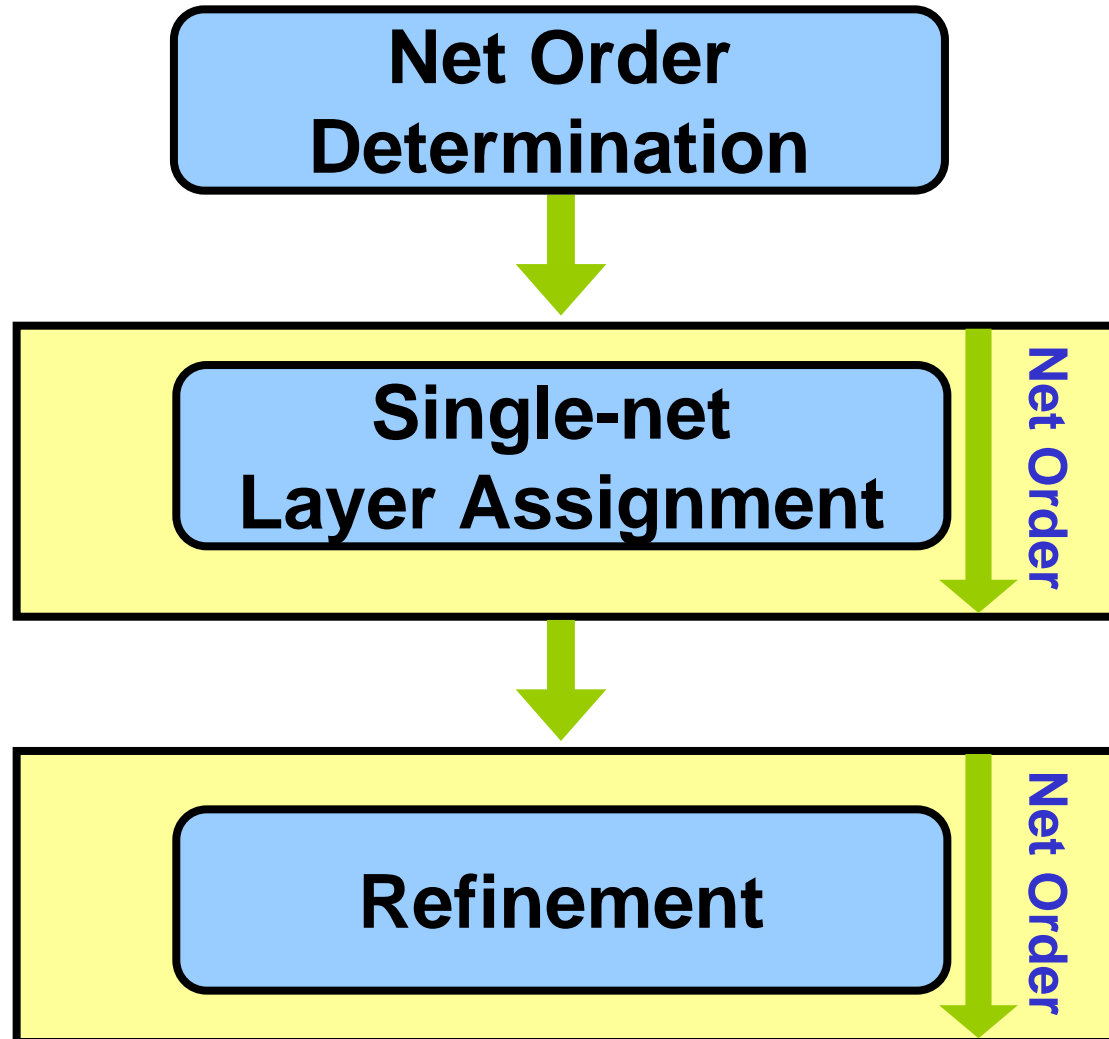
Problem Formulation

- Given a 2D routing result, finds a 3D counterpart through layer assignment
 - Minimize **via overflow**, and **wirelength**
 - Keep the same **wire overflow** from 2D routing result



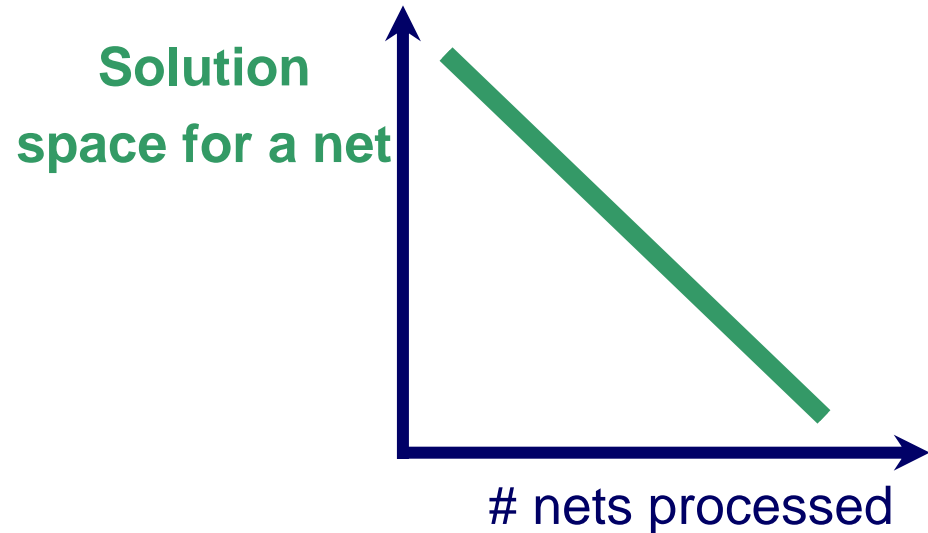
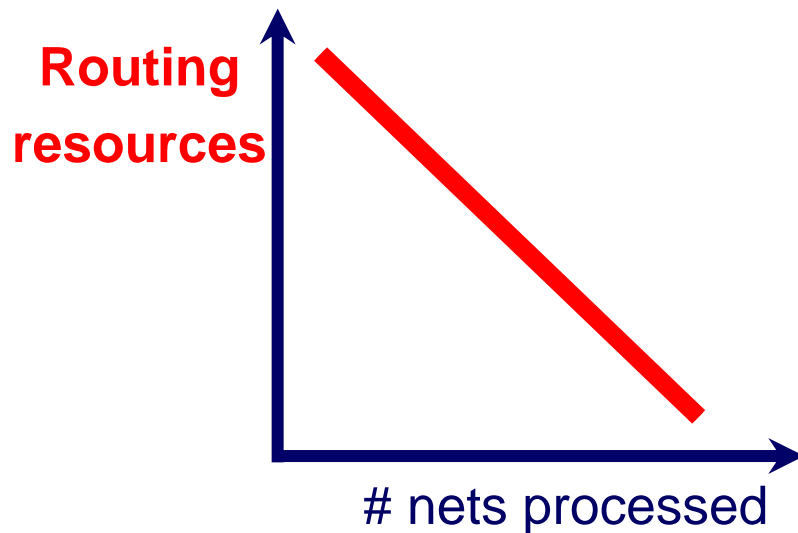
Algorithm

- Our algorithm contains 3 steps



Net Order Determination

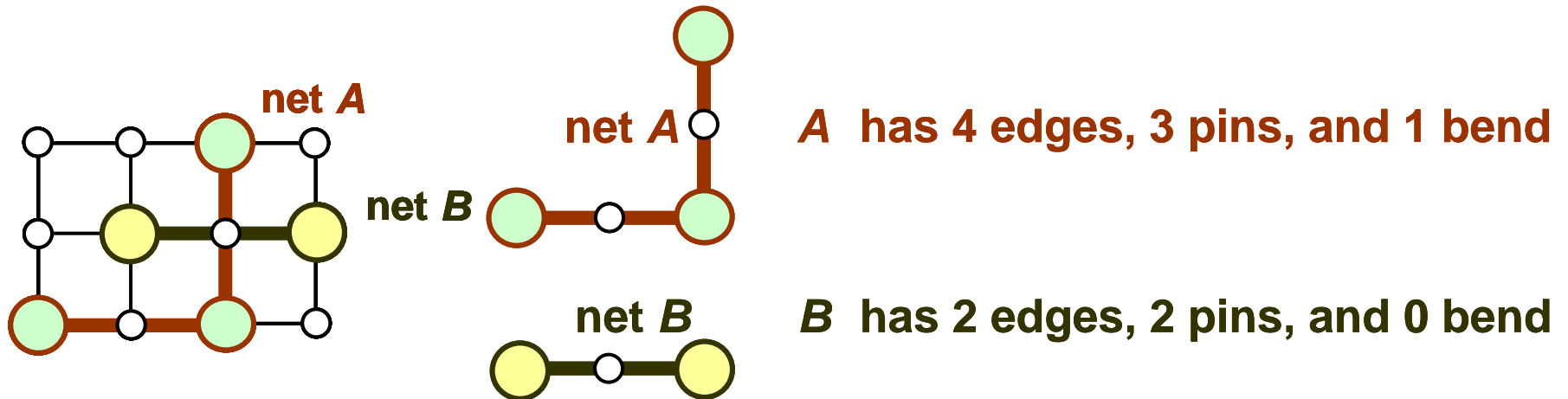
- Due to the **limited routing resources**, a net processed earlier has **larger solution space**



- **Net order** should maximize resource utilization

Net Order Determination

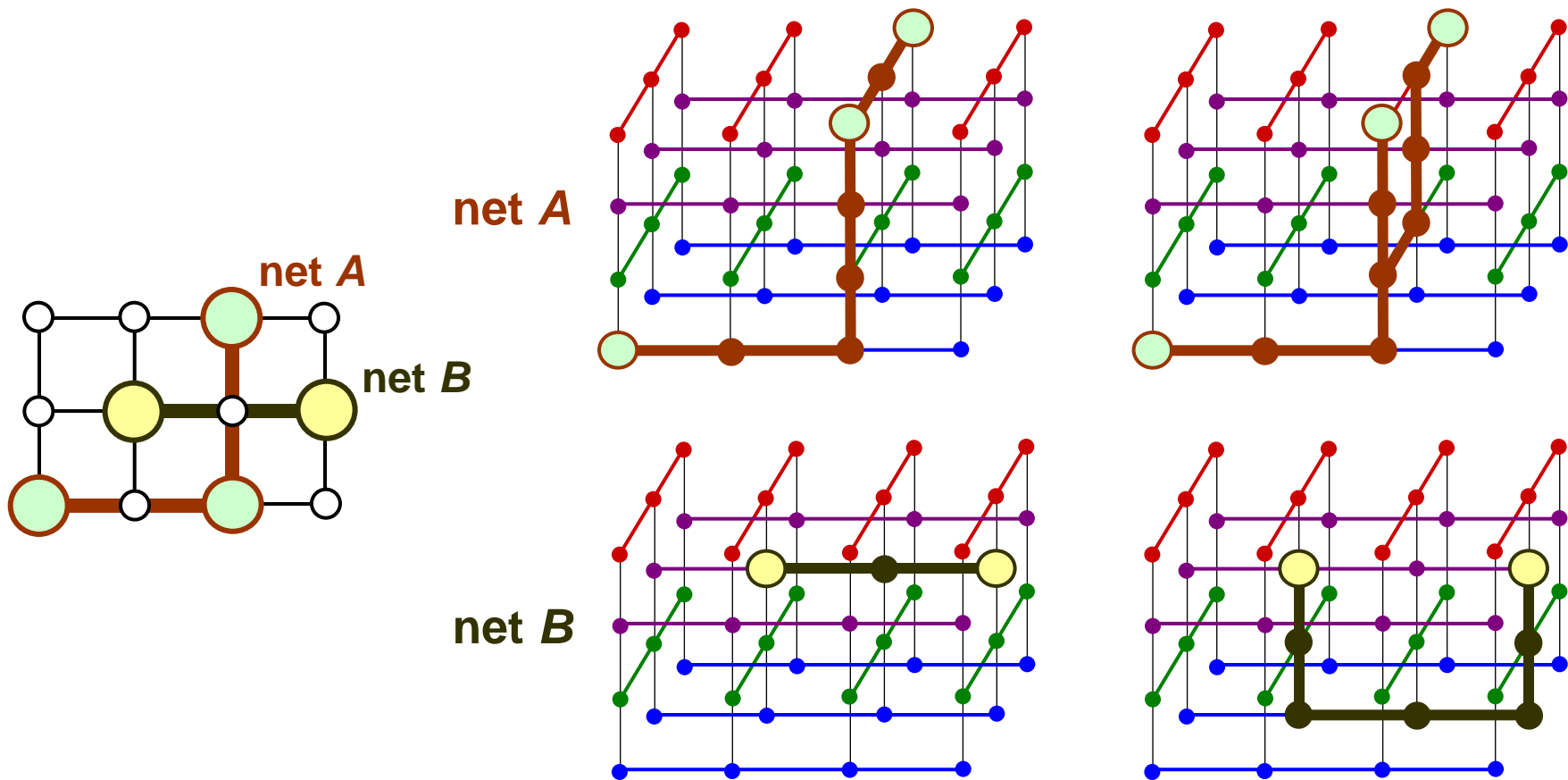
- For each 2D net T , we use 3 parameters to determine its order — **Score(T)**
 - **Length(T)** : # of edges in T
 - **PinNum(T)** : # of pins in T
 - **Bends(T)** : # of bends in T
- **Net Order** derived from sorting $Score(T)$ for each net T decreasingly



Net Order Determination — *Length*

- A net with **longer length** will **occupy more routing resources**

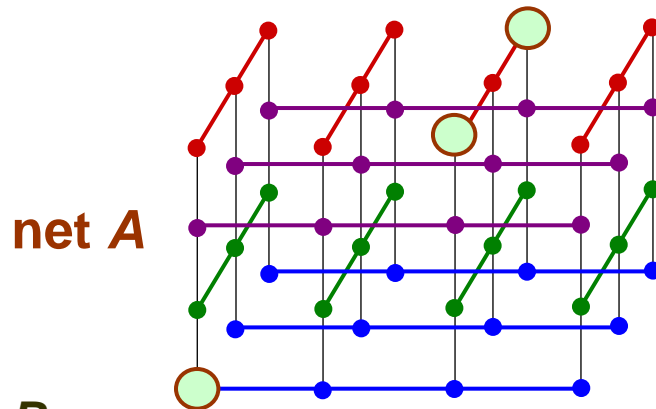
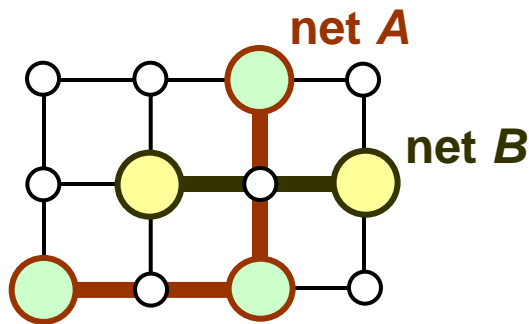
– $Length(T) \uparrow$, $Score(T) \downarrow$



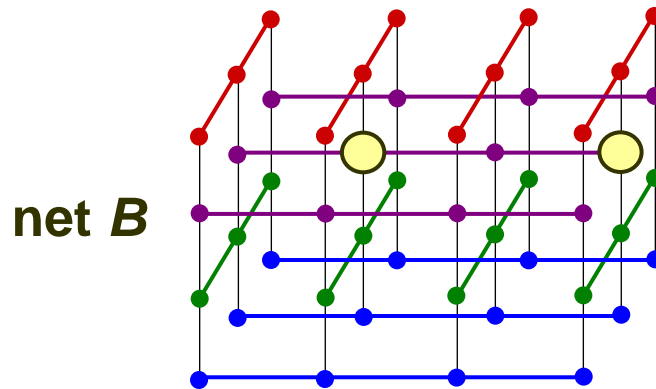
Net Order Determination — *PinNum*

- The role of **pin** in global routing, just like the role of **checkpoints** in race

– $PinNum(T) \uparrow$, $Score(T) \uparrow$



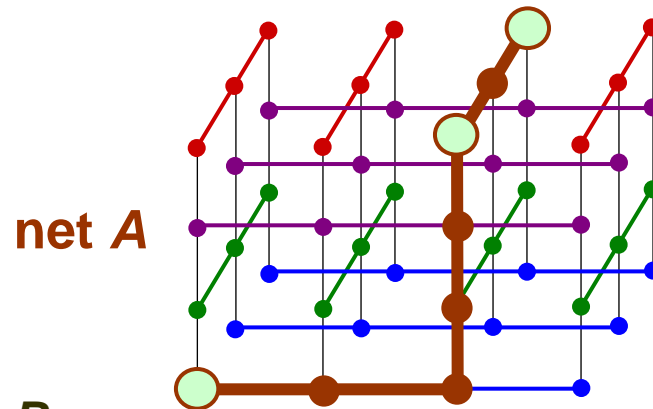
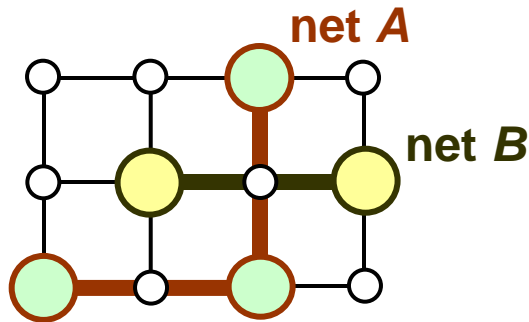
3 pins = 3 checkpoints



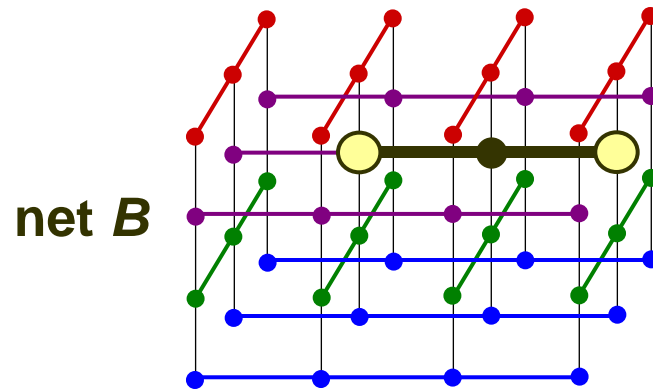
2 pins = 2 checkpoints

Net Order Determination — *Bends*

- Changing routing direction needs vias, so bend is similar to pin
 - $Bends(T) \uparrow$, $Score(T) \uparrow$



1 bend at least needs an via



0 bend may not need vias

Net Order Determination

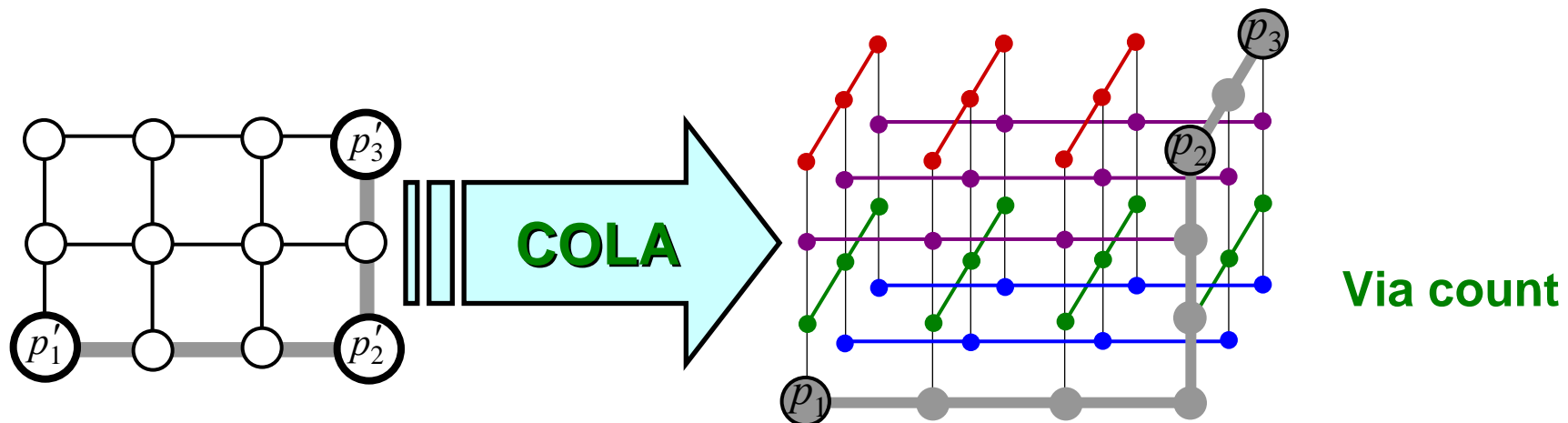
- $Score(T)$'s relationships with $Length(T)$, $PinNum(T)$, and $Bends(T)$.
 - $Length(T) \uparrow$, $Score(T) \downarrow$
 - $PinNum(T) \uparrow$, $Score(T) \uparrow$
 - $Bends(T) \uparrow$, $Score(T) \uparrow$
- **$Score(T)$**
 - = $(\alpha \times Bends(T) + \beta \times PinNum(T)) / Length(T)$

Single-net Layer Assignment

□ [Lee et al., TCAD'08]

“Congestion-Constrained Layer Assignment for Via Minimization in Global Routing”

- COLA finds a layer assignment result with minimum via count for a 2D net



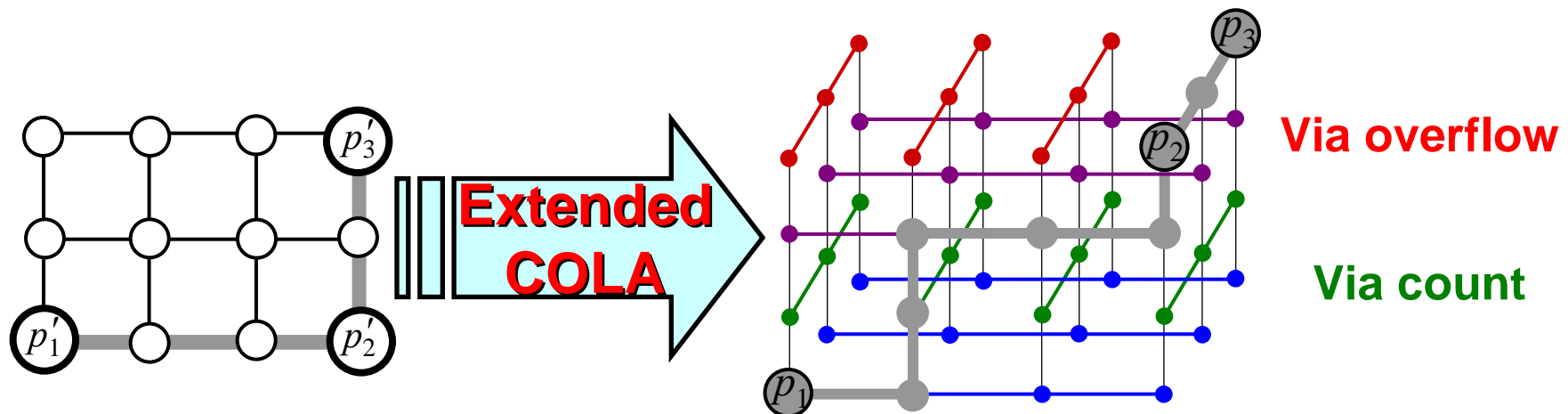
Single-net Layer Assignment

- [Lee et al., TCAD'08]

- “*Congestion-Constrained Layer Assignment for Via Minimization in Global Routing*”

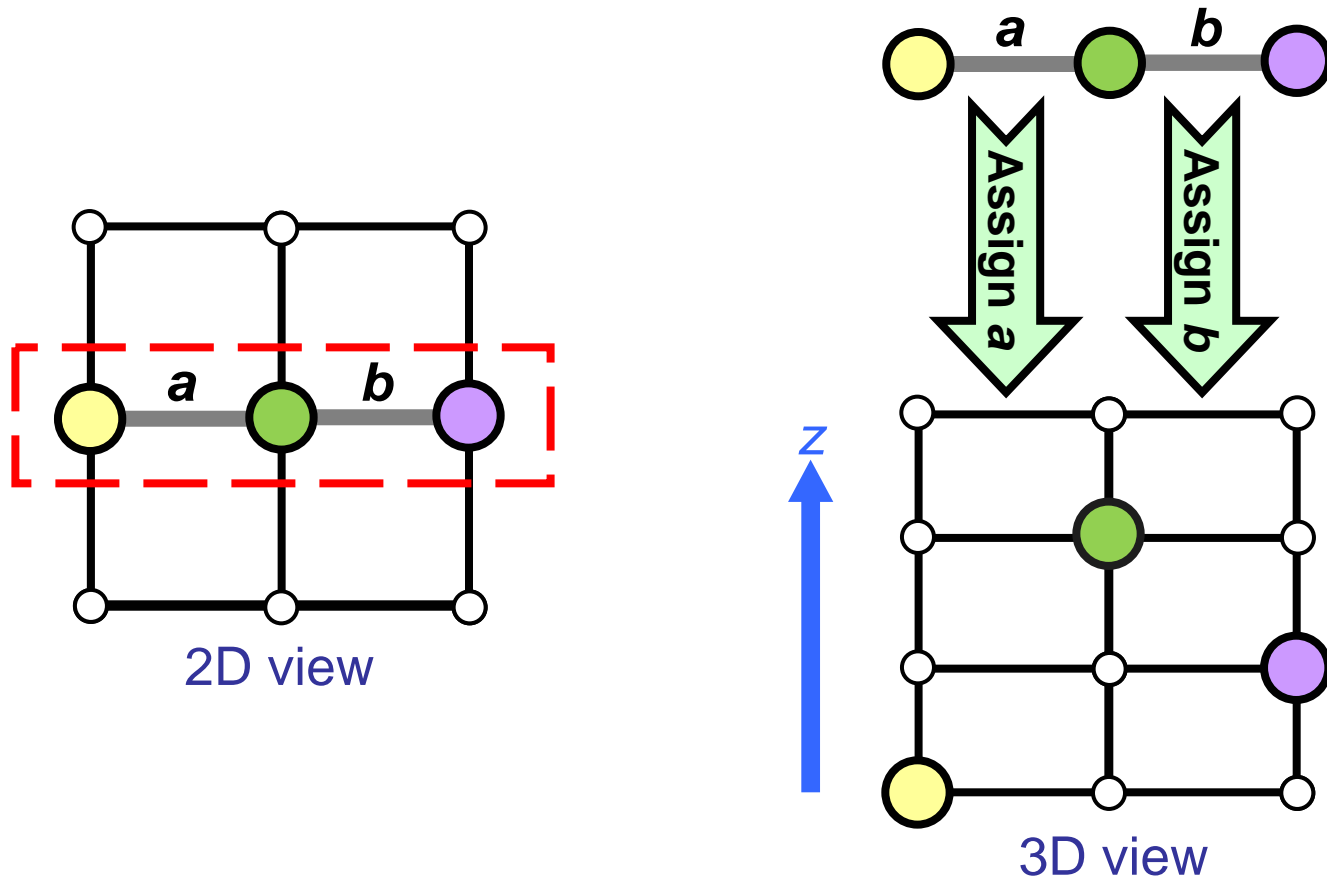
- COLA finds a layer assignment result with minimum via count for a 2D net

- Extends from COLA, our algorithm can deal with via overflow and via count



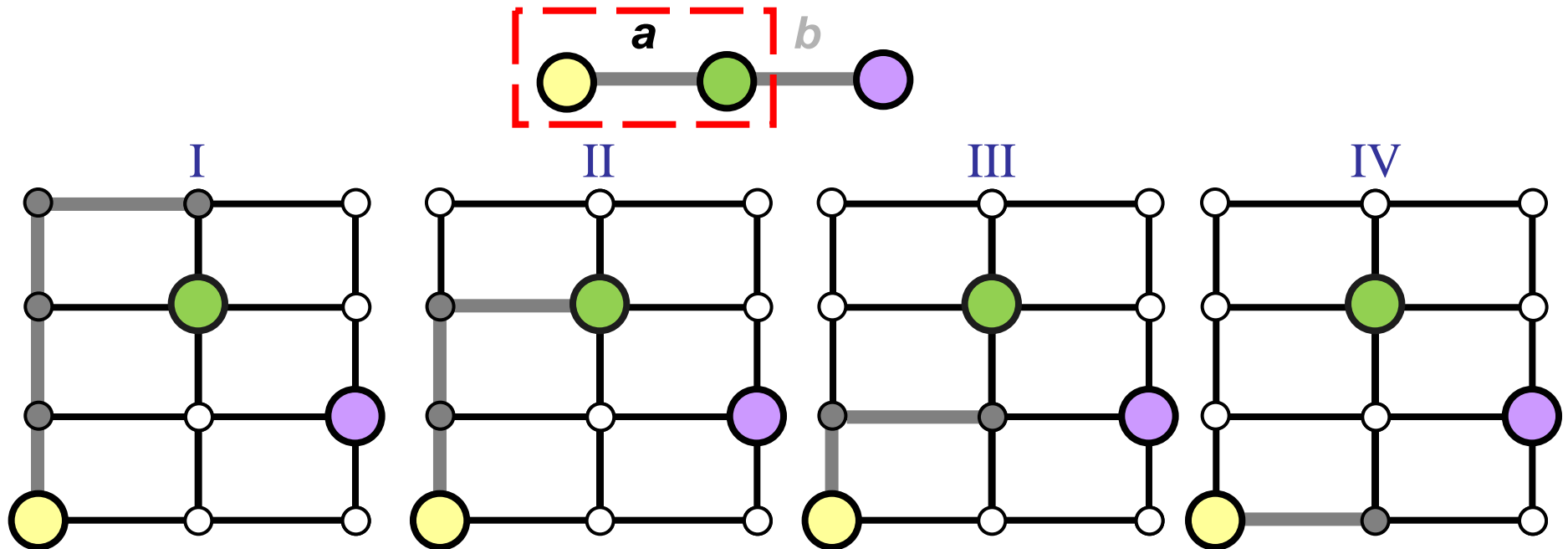
Single-net Layer Assignment

- DP-based layer assignment method
 - Minimize increase on **via overflow**, and **via count**
- For a net, assign one edge at a time



Single-net Layer Assignment

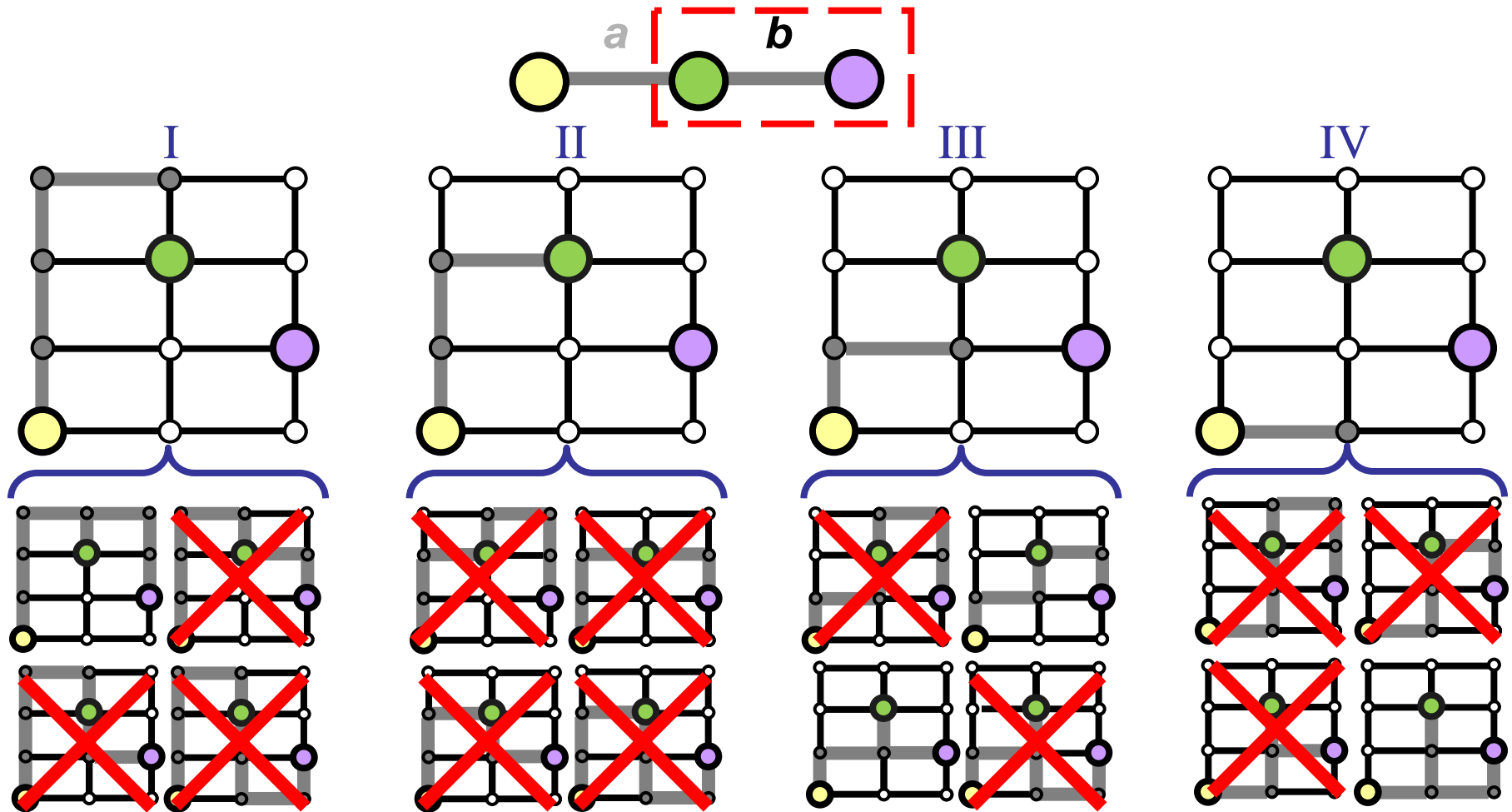
- Vias are placed after edges are assigned
 - Vias are determined by edges and pins



- Since vias on different tiles are independent, the via overflow increase on each tile can be calculated independently

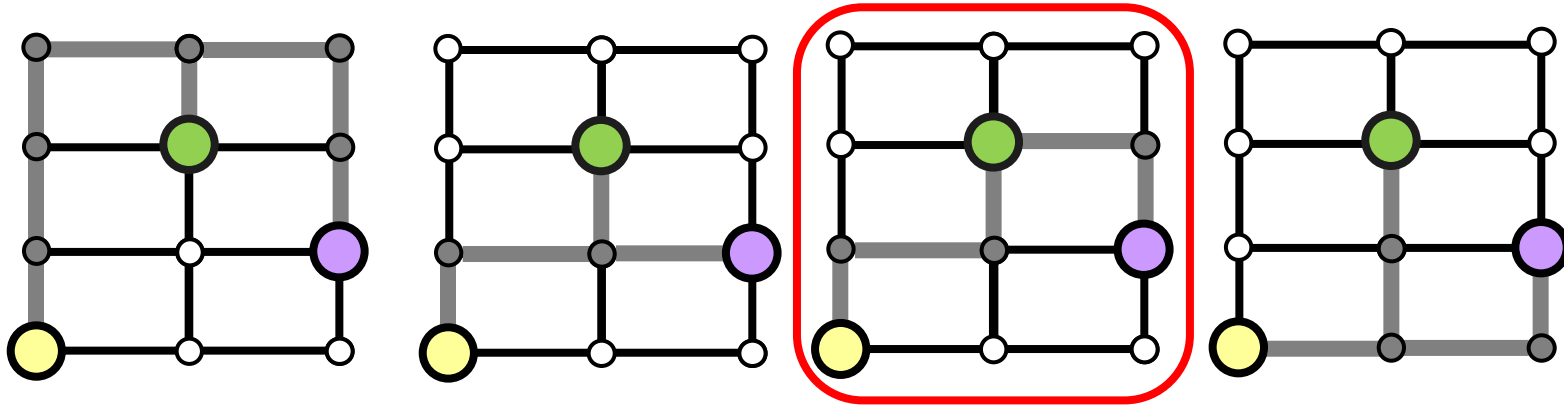
Single-net Layer Assignment

- Memorize the minimum via overflow one on each stage and propagate it to the next stage



Single-net Layer Assignment

- In the end, the optimal result is the one with **minimum total via overflow increase**

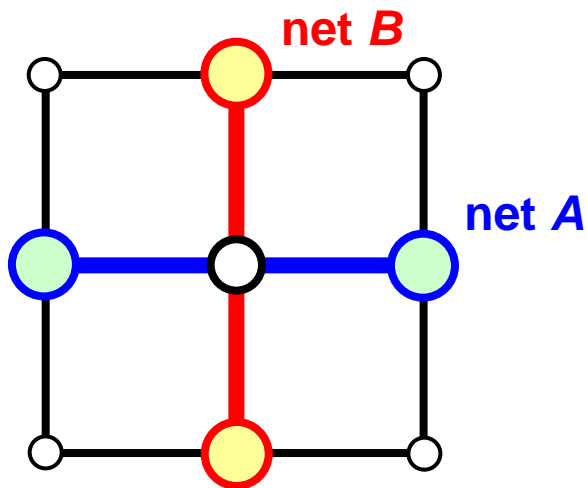


Optimal result

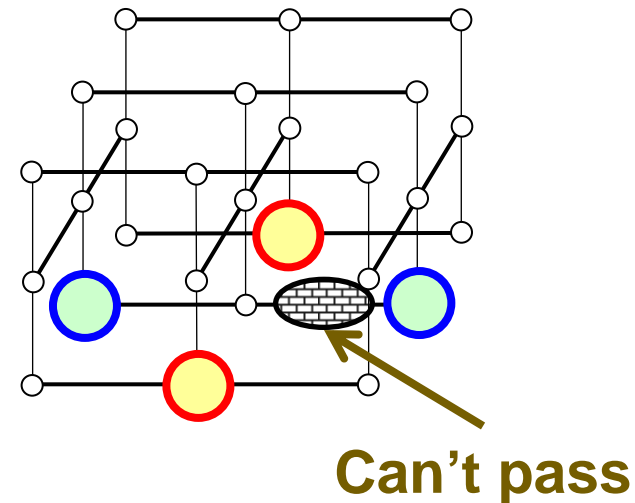
- If there is a tie on via overflow increase, choose the one with **minimum via count**

Refinement

- ❑ Refinement can improve the via overflow of the original layer assignment result
 - Rip-up and re-layer assignment for each net
- ❑ Enhance the net order



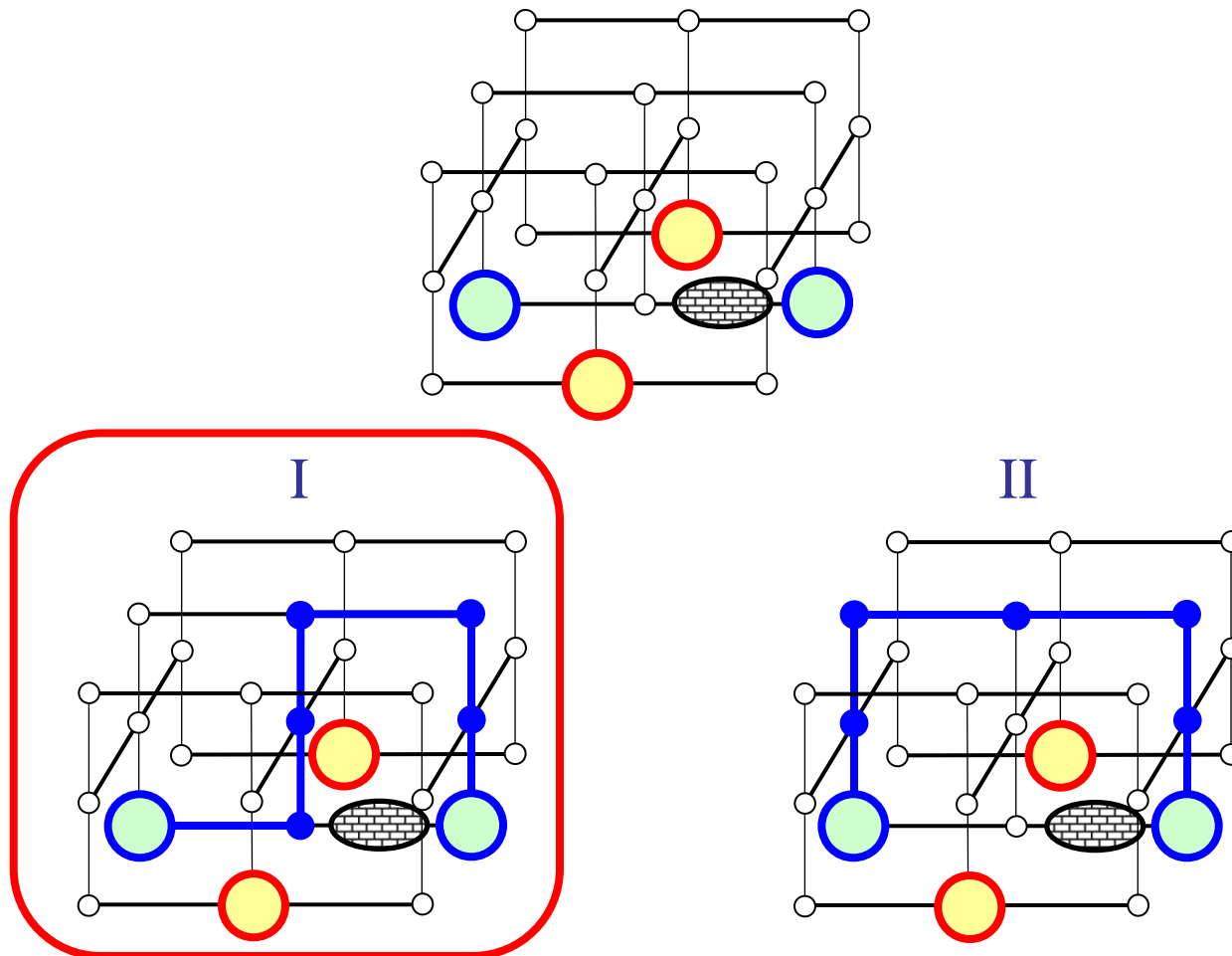
2D view



3D view

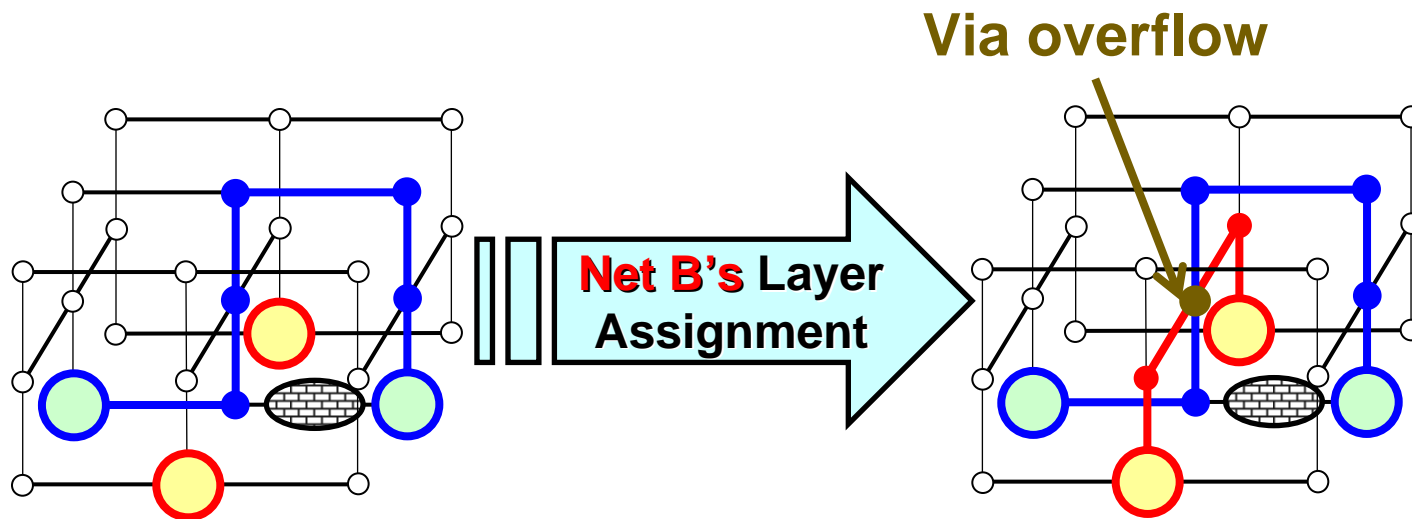
Refinement

- Net A assigned first, and net B assigned second
 - Net A has 2 choices with the same via overflow increase and via count

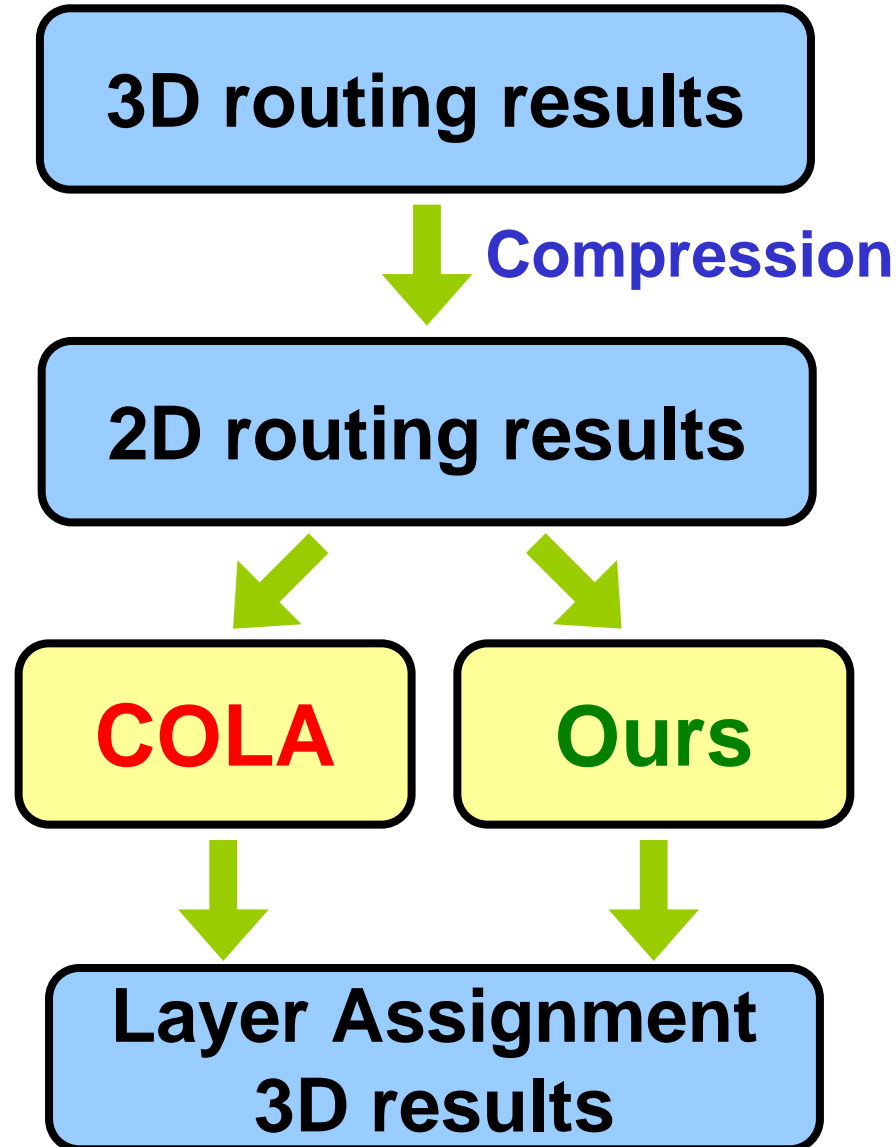


Refinement

- If net *A* chooses improperly, net *B* will generate via overflow inevitably
 - It is impossible for net *A* to know how to choose before net *B* is assigned

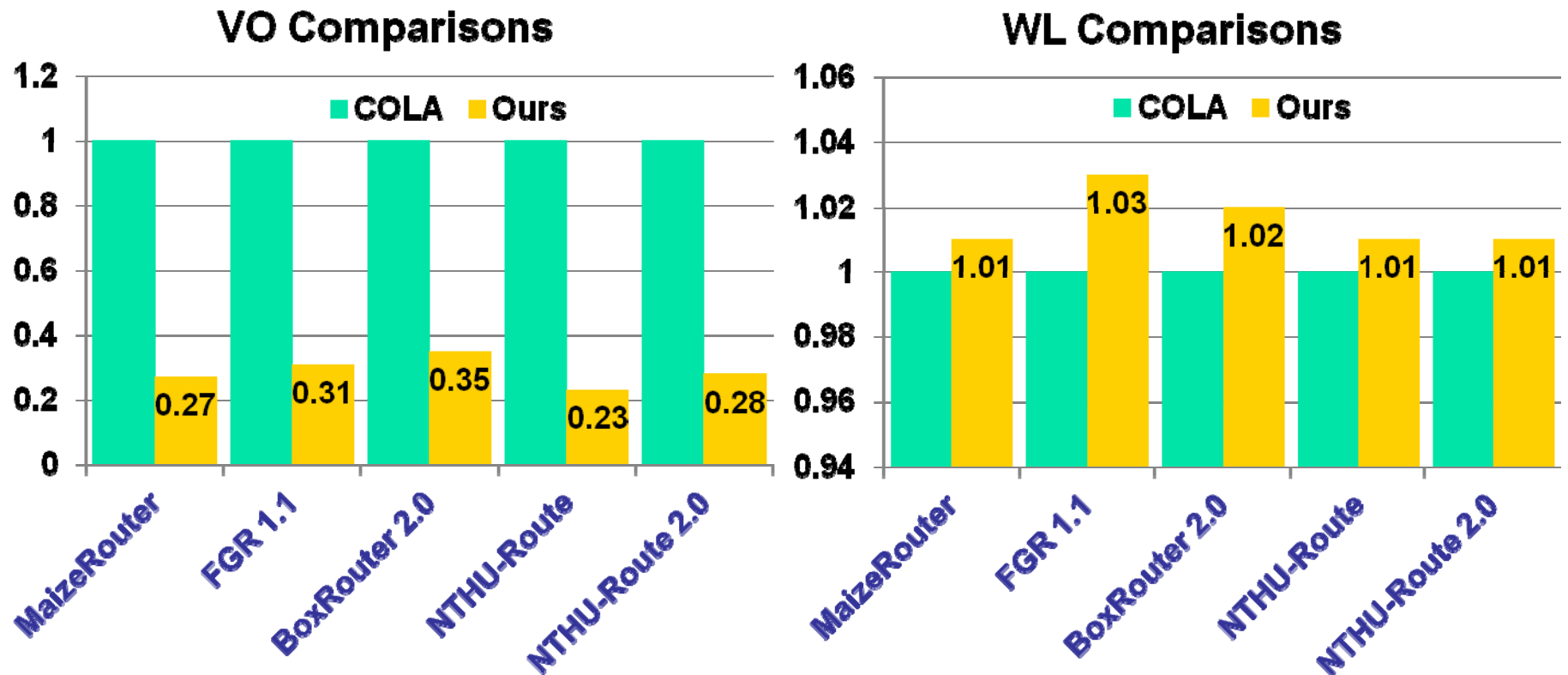


Experiment Flow



Experimental Results without Refinement

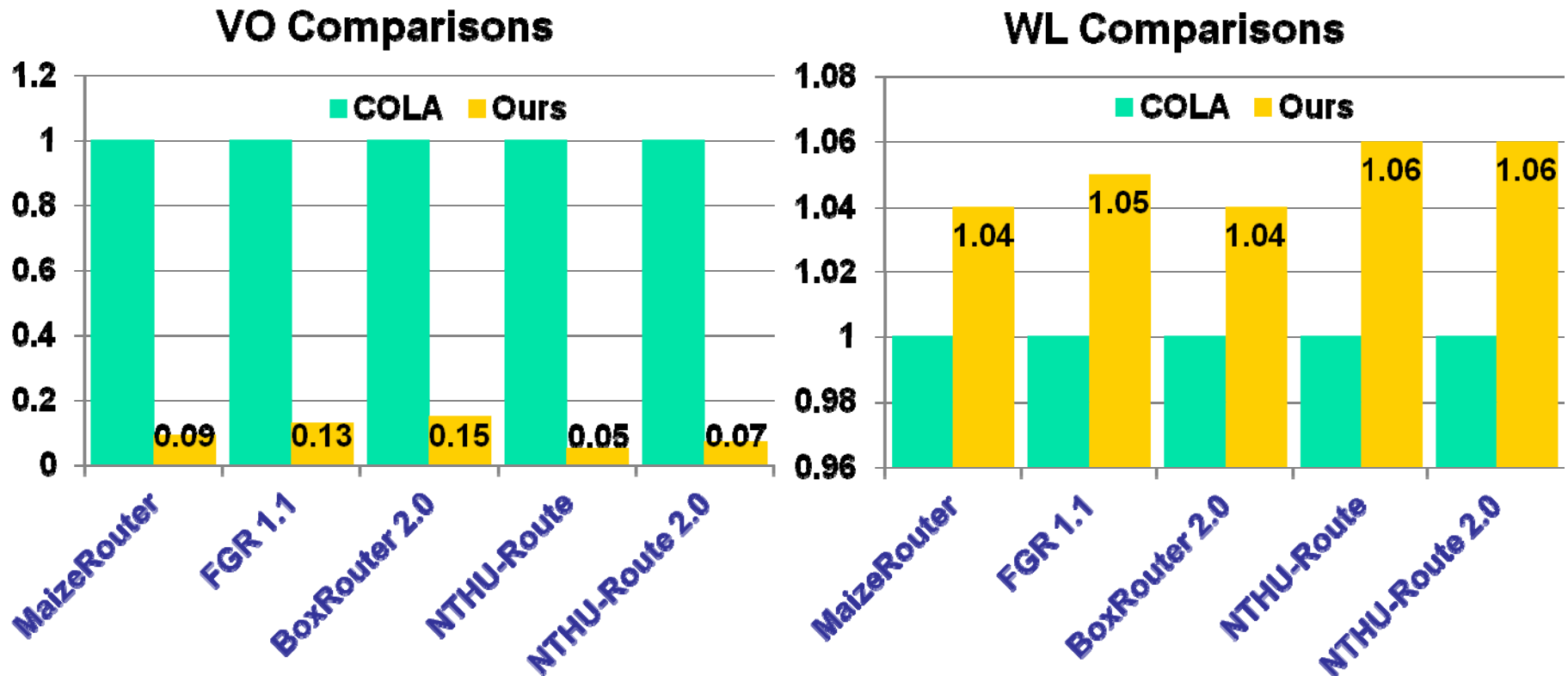
- Without refinement, our algorithm induced **23~35%** via overflow with **1~3%** WL increase compared with COLA [1]



[1] Lee et al., TCAD'08, "Congestion-Constrained Layer Assignment for Via Minimization in Global Routing"

Experimental Results with Refinement

- With refinement, our algorithm induced 5~15% via overflow with 4~6% WL increase compared with COLA [1]



[1] Lee et al., TCAD'08, "Congestion-Constrained Layer Assignment for Via Minimization in Global Routing"

Conclusions

- ❑ Develop a layer assignment algorithm considering **via overflow** and **via count**
- ❑ **Future work:** a more effective **net order** and **layer assignment method**

Q & A

Thank You!
and

